

Tussle in Domain Namespace



donderdag 2 mei 2019
SNOW ❄️

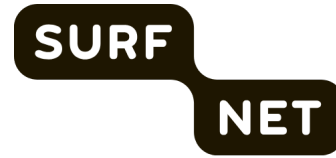


Wat is/Wat doet NLNETLABS

- Non-profit stichting – sinds 1999 – subsidies & donaties



VERISIGN



Wat is/Wat doet NLNETLABS

- Missie:

Leveren van globaal erkende innovaties en expertise in die technologieën die een netwerk van netwerken maken tot een Open Internet voor allen.

- Doel:

- *Ontwikkelen van Open Source software en Open Standaarden ten behoeve van het Internet.*

Wat is/Wat doet NLNETLABS

- Doel:
 - Ontwikkelen van *Open Source software* en *Open Standaarden* ten behoeve van het Internet.



NSD



unbound



ROUTINATOR

Krill

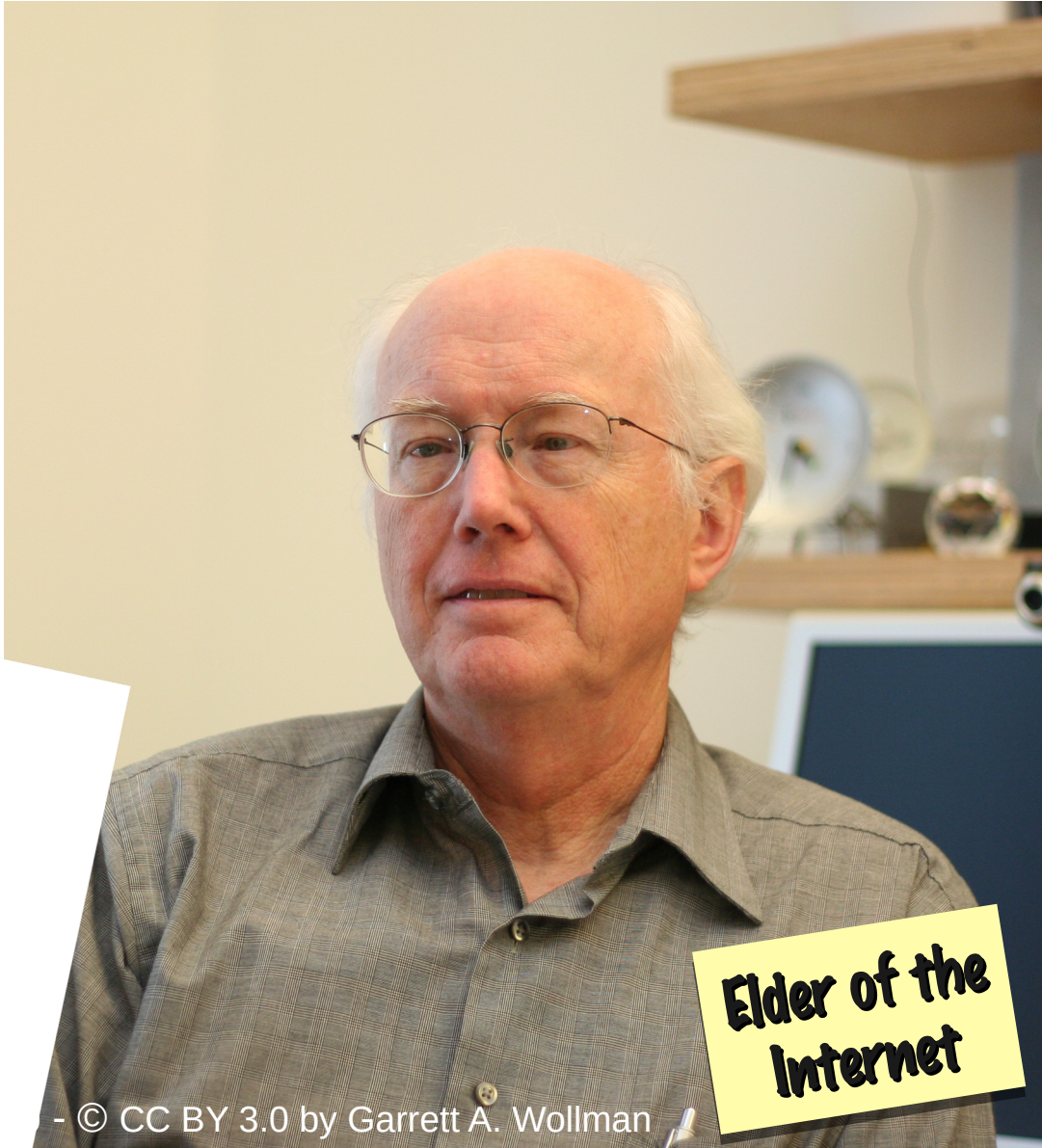
- Idns
- Net::DNS
- Net::DNS::SEC

Research – Internet metingen – Studenten projecten

Tussle

bakkeleien ; plukharen

<https://www.mijnwoordenboek.nl/vertaal/NL/EN/tussle>



Elder of the Internet

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Tussle in Cyberspace: Defining Tomorrow's Internet

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Abstract

The architecture of the Internet is based on a number of principles, including the self-describing datagram packet, the end to end arguments, diversity in technology and global addressing. As the Internet has moved from a research curiosity to a recognized component of mainstream society, new requirements have emerged that suggest new design principles, and perhaps suggest that we revisit some old ones. This paper explores one important reality that surrounds the Internet today: different stakeholders that surround the Internet milieu have interests that are not shared with each other, and these interests are often in conflict.

1. INTRODUCTION

The Internet was created in simpler times. Its creators and early users shared a common goal—they wanted to build a network infrastructure to hook all the computers in the world together so that as yet unknown applications could be invented to run there. All the players—users, operators or operators, shared a common sense of purpose.

Tussle

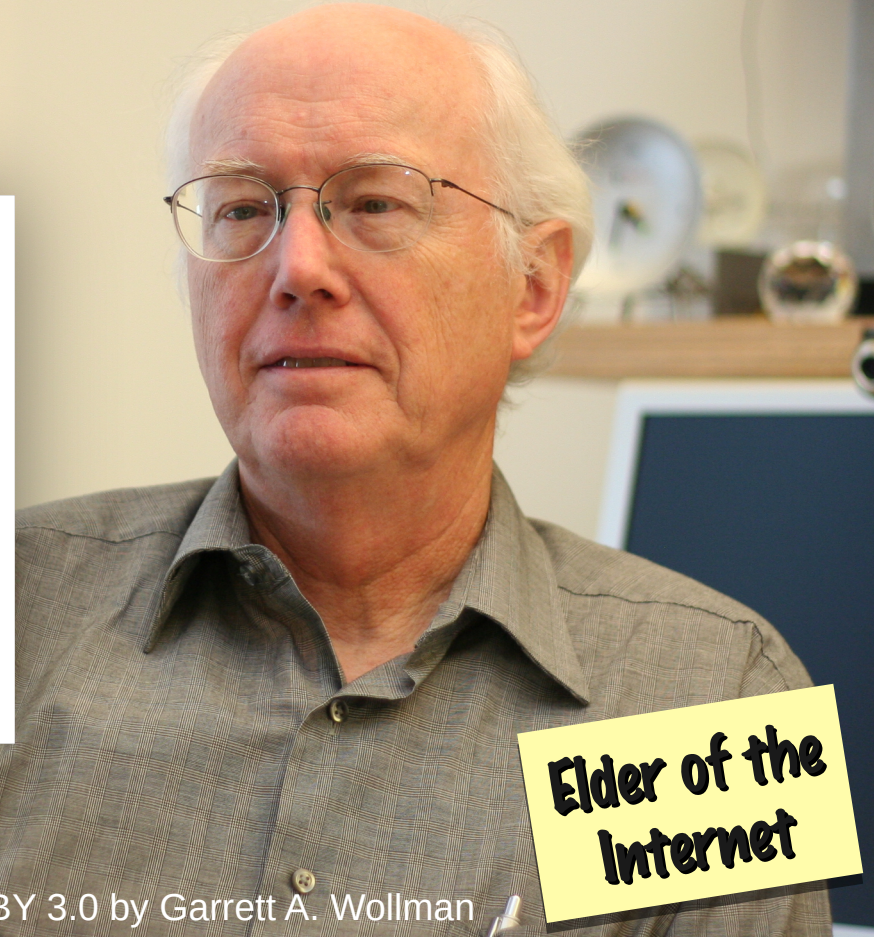
bakkeleien ; plukharen

<https://www.mijnwoordenboek.nl/vertaal/NL/EN/tussle>

2. PRINCIPLES

In this paper we offer some design principles to deal with tussle. Our highest-level principle is:

- Design for variation in outcome, so that the outcome can be different in different places, and the tussle takes place within the design, not by distorting or violating it. Do not design so as to dictate the outcome. Rigid designs will be broken; designs that permit variation will flex under pressure and survive.



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Tussle

bakkeleien ; plukharen

<https://www.mijnwoordenboek.nl/vertaal/NL/EN/tussle>

2. PRINCIPLES

Within this guiding principle, we identify two more specific principles:

- Modularize the design along tussle boundaries, so that one tussle does not spill over and distort unrelated issues.
- Design for choice, to permit the different players to express their preferences.

Abstract

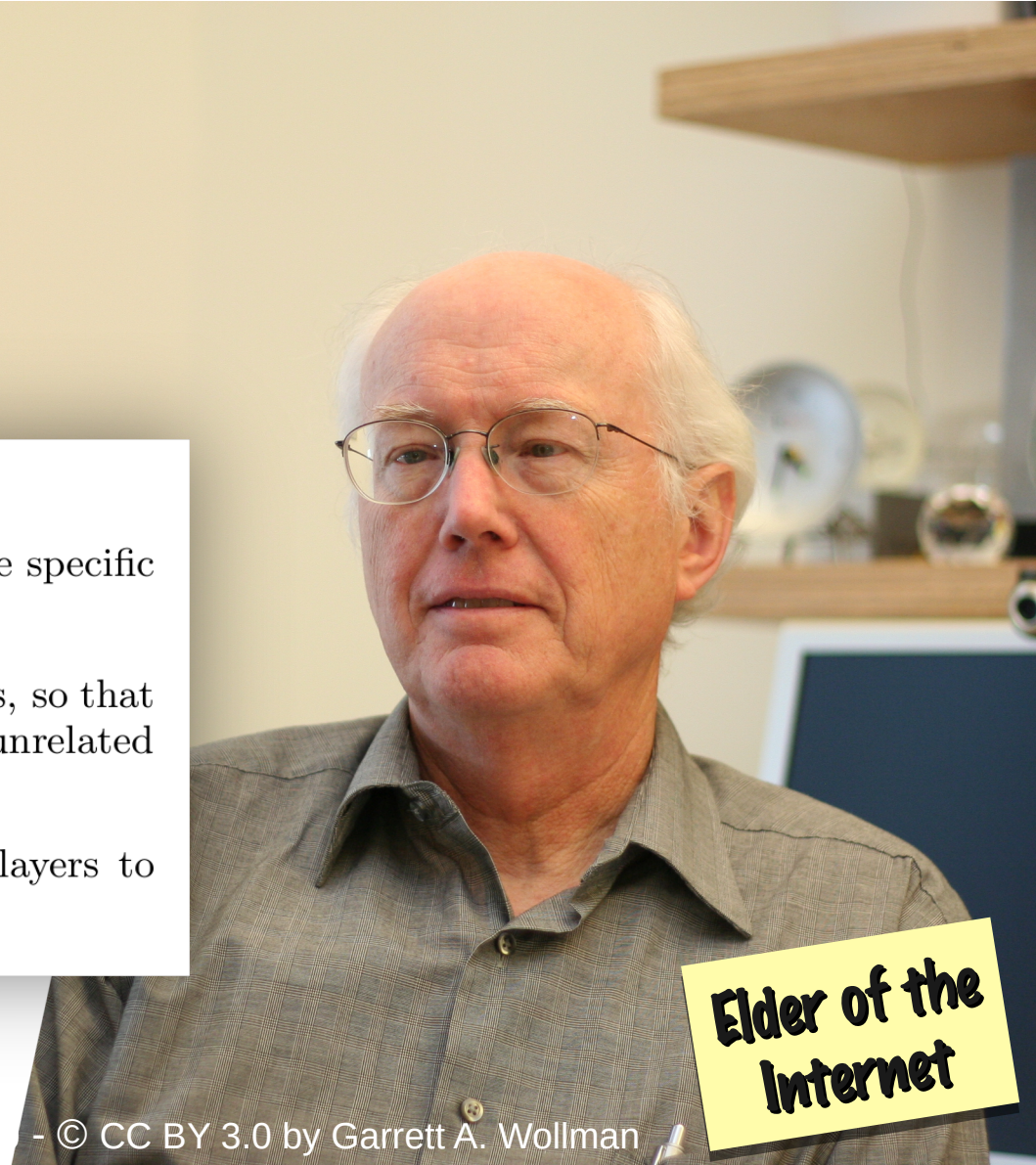
The architecture principles, including the end-to-end arguments, diversity in technology and global connectivity to a recognized component of mainstream society, and perhaps requirements have emerged that suggest new design principles, and this paper explores one important reality that underlies the Internet today: different stakeholders that surround each other, and these stakeholders have interests that are not necessarily aligned.

1. INTRODUCTION

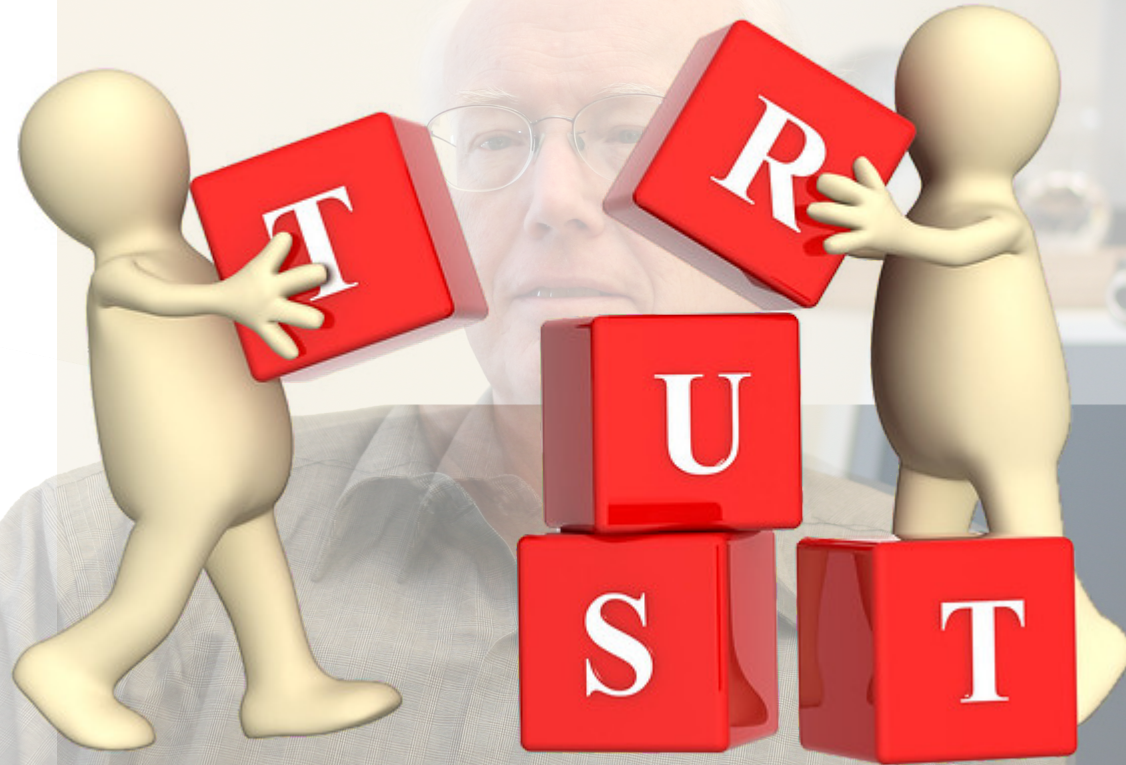
The Internet was created in simpler times. Its creators and early users shared a common goal—they wanted to build a network infrastructure to hook all the computers in the world together so that as yet unknown applications could be invented to run there. All the players, whether users or operators, shared a common sense of purpose.

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Elder of the Internet



Tussle Spaces



Tu

Internet

Information Sciences Institute
braden@isi.edu

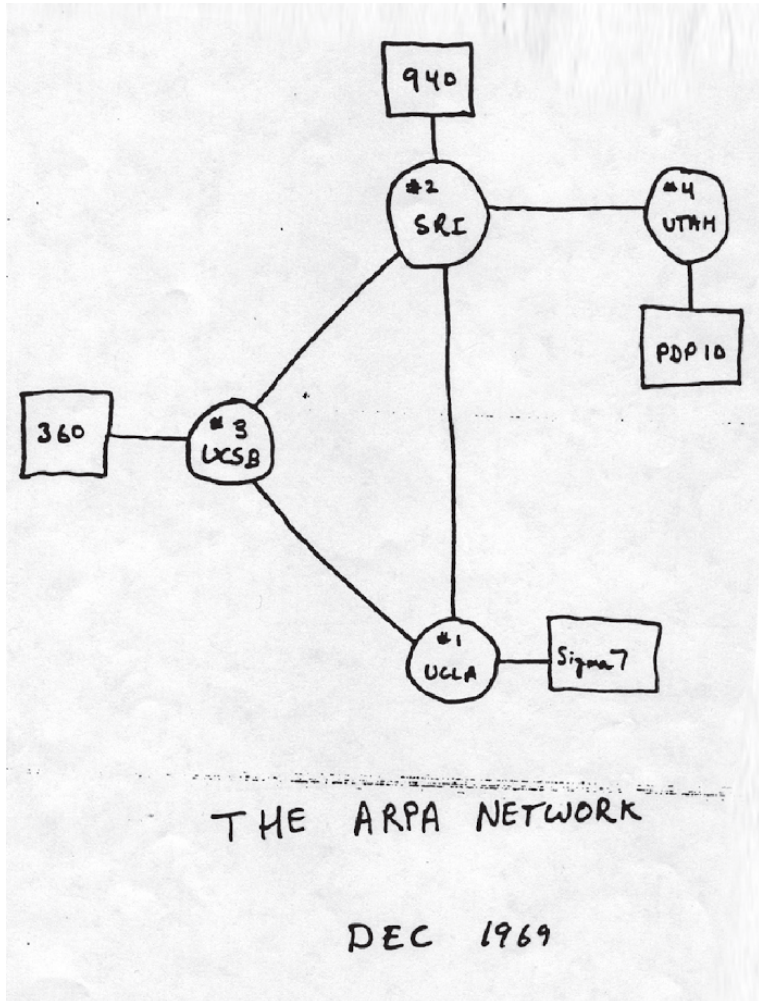
1. INTRODUCTION

The Internet was created in simpler times. Its creators and early users shared a common goal—they wanted to build a network infrastructure to hook all the computers in the world together so that as yet unknown applications could be invented to run there. All the players—users or operators, shared a sense of purpose.

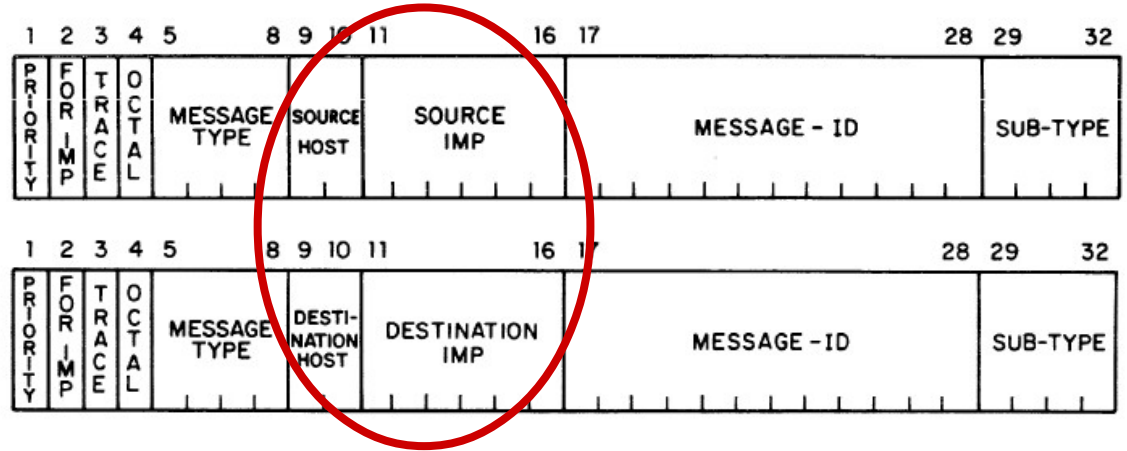
Abstract

The architecture of the Internet is based on a number of principles, including end-to-end argument, scalability, and simplicity. As the Internet has moved from a research curiosity to a recognized component of mainstream society, new requirements have emerged that suggest new design principles, and perhaps suggest that we revisit some old ones. This paper explores one important reality that suggests the Internet today: different stakeholders that surround the Internet milieu have interests that are not necessarily aligned with each other, and these interests are often in conflict.

Namespace op het internet



NCP (Network Control Program)



- December 1973
HOSTS.TXT (RFC 606)

Namespace op het internet

NCP (Network Control Program)

ARPANET DIRECTORY
NIC 19275
Jan. 1974

HOST NAMES

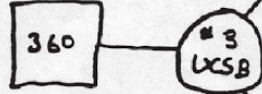
HOST NAMES

HOSTNAME	HOST ADDR (Dec)	LIAISON	STATUS
AFWL-TIP	176	D Hyde (505)247-1711 x3803	TIP, Up 3-74
ALOHA-TIP	164	R Binder (808)948-7066	TIP
AMES-11	208	J Hart (415)965-5935	USER, up 12-73
AMES-67	16	W Hathaway (415)965-6033	SERVER
AMES-TIP	144	W Hathaway (415)965-6033	TIP
ANL	?	L Amiot (312)739-7711 x4309	SERVER, up 2-74
ARPA-DMS	28	S Crocker (202)694-5037	USER, Agency use only
ARPA-TIP	156	S Crocker (202)694-5037	TIP
BBN-11X	5	R Thomas (617)491-1850 x483	Peripheral processor for #69, up 12-73
BBN-1D	232	A McKenzie (617)491-1850 x441	USER
BBN-NCC	40	A McKenzie (617)491-1850 x441	USER
BBN-TENEX	69	R Thomas (617)491-1850 x483	SERVER
BBN-TENEXB	133	R Thomas (617)491-1850 x483	SERVER, Limited
BBN-TESTIP	158	A McKenzie (617)491-1850 x441	TIP (magtape)
BELVOIR	27	W Andrews (703)664-5511	USER, up 6-74
BRL	29	M Romanelli (301)278-4574	USER
CASE-10	13	J Calvin (216)368-2984	SERVER
CCA-TENEX	31	R Winter (617)491-3670	SERVER
CCA-TIP	159	R Winter (617)491-3670	TIP
CMU-10A	78	H Van Zoeren (412)621-2600 x160	SERVER

28	29	32
- ID	SUB-TYPE	

28	29	32
-ID	SUB-TYPE	

C 606)



THE ARPANET

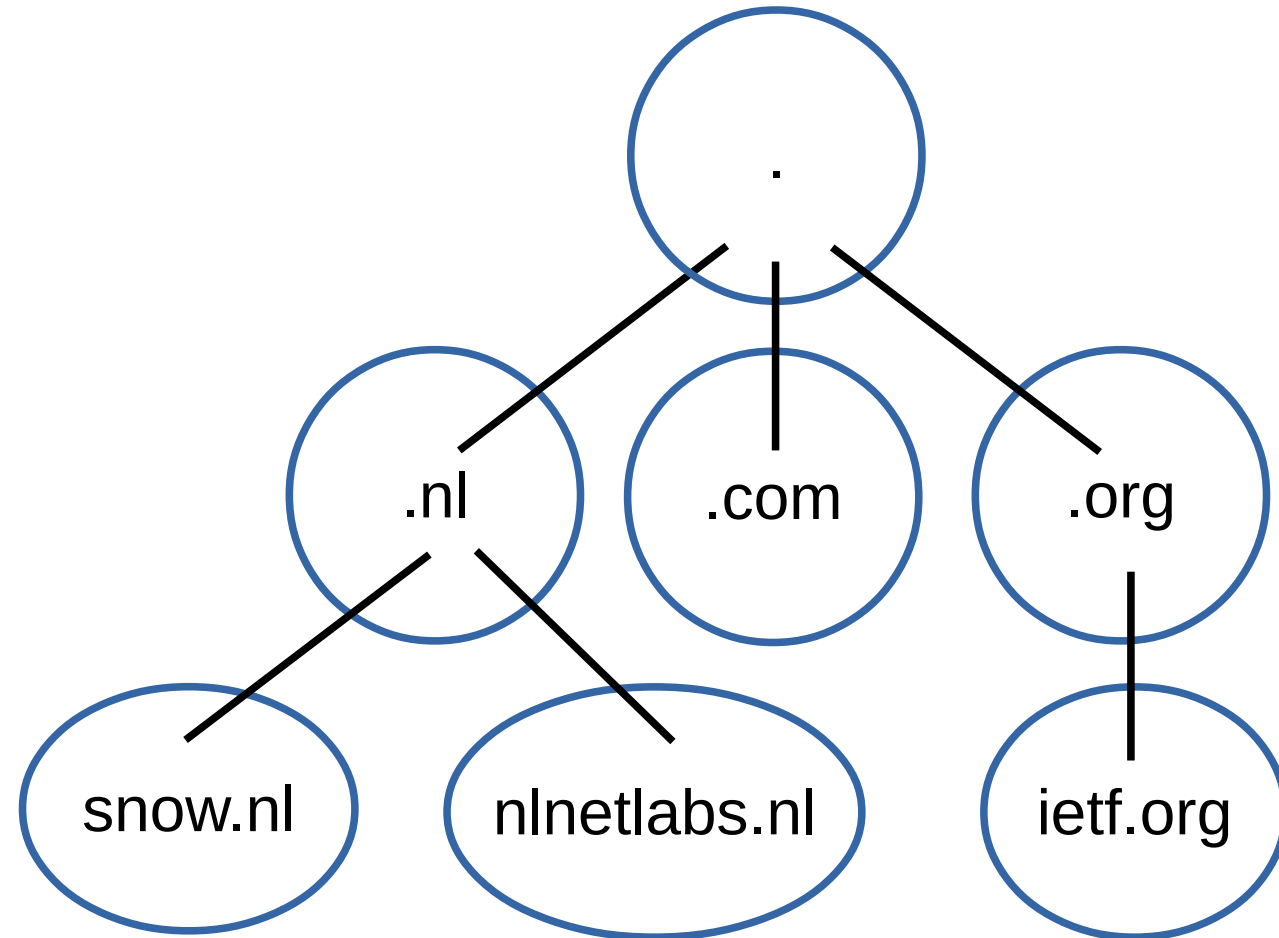
DEC

Namespace op het internet



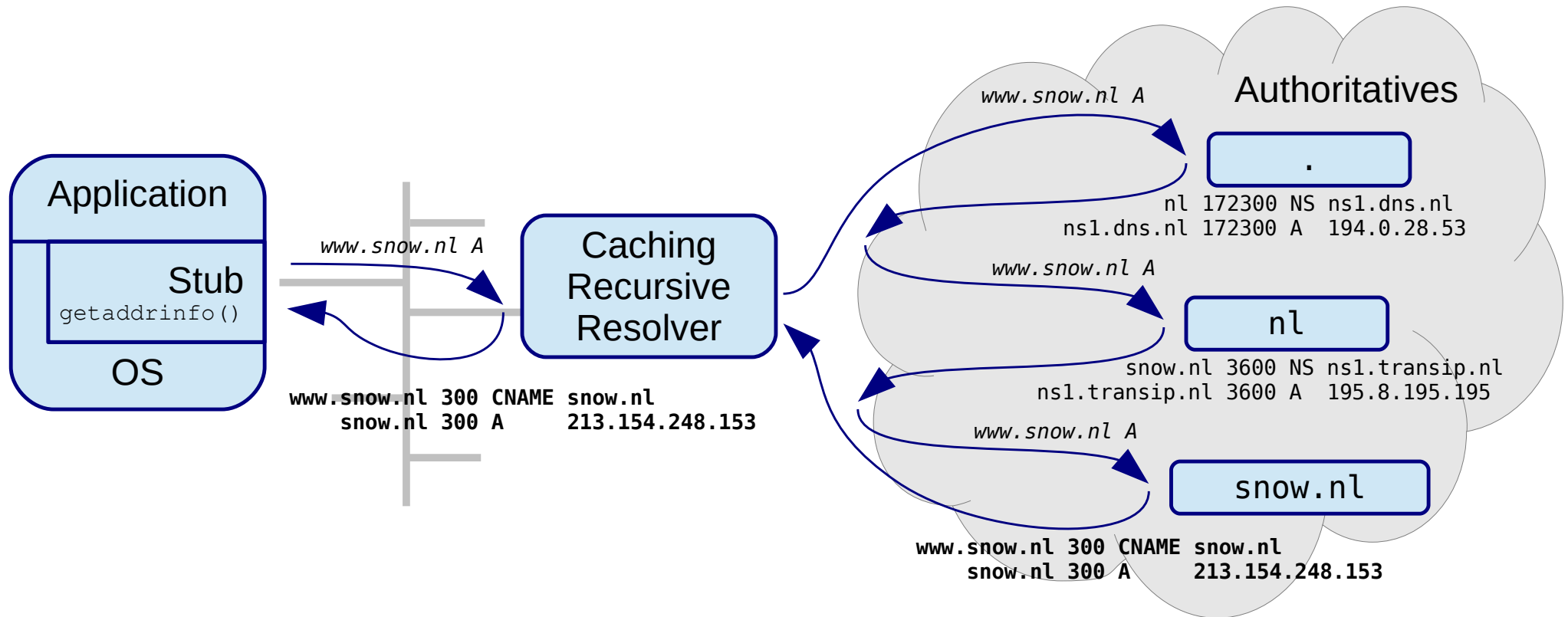
- 1 januari 1983 NCP → IP/TCP
flagday
- max 256 → max 4.294.967.296 hosts
- november 1983 DNS (RFC 882)
Domain Name System
- november 1987 STD13
(RFC 1034 & RFC 1035)
- Niet alleen IP adressen (ook mail)

Domain Namespace - schaal



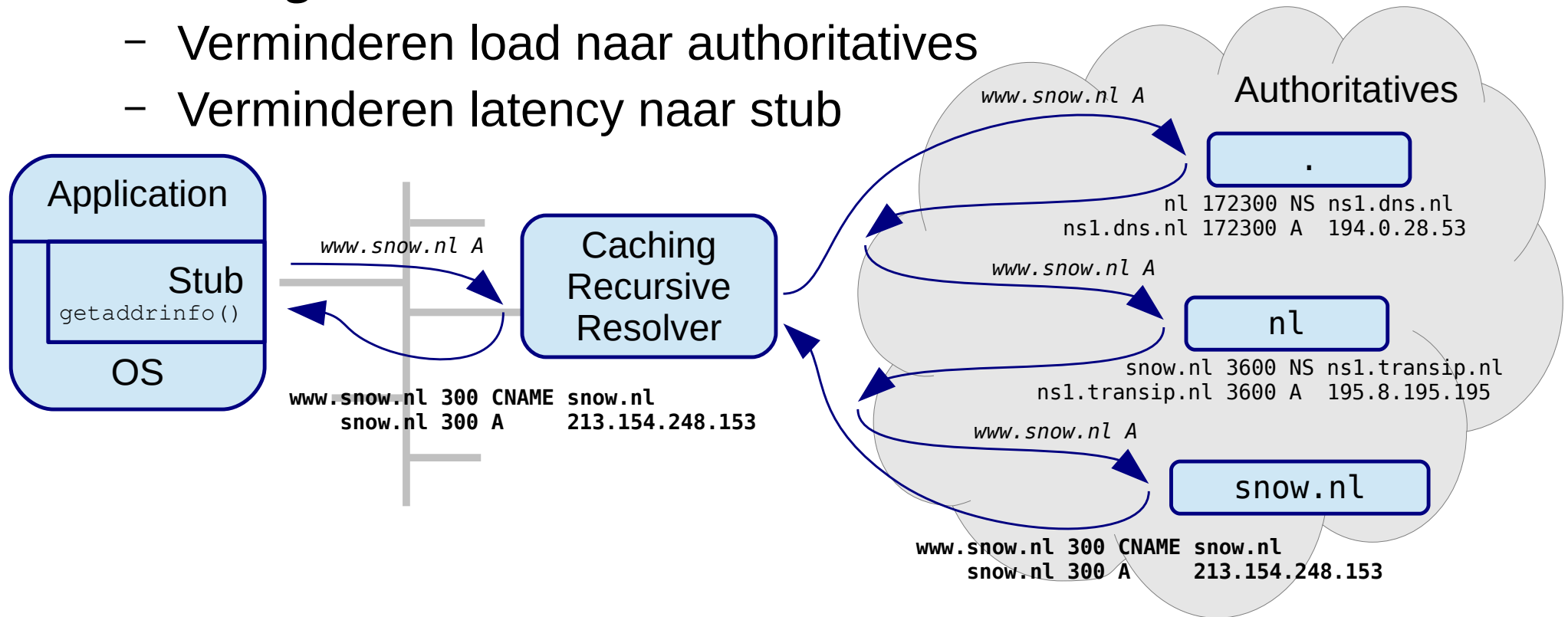
- 13 root servers in 12 organisaties
- 1532 tlds
- .com 141.000.000
.nl 5.000.000

Domain Name System - schaal



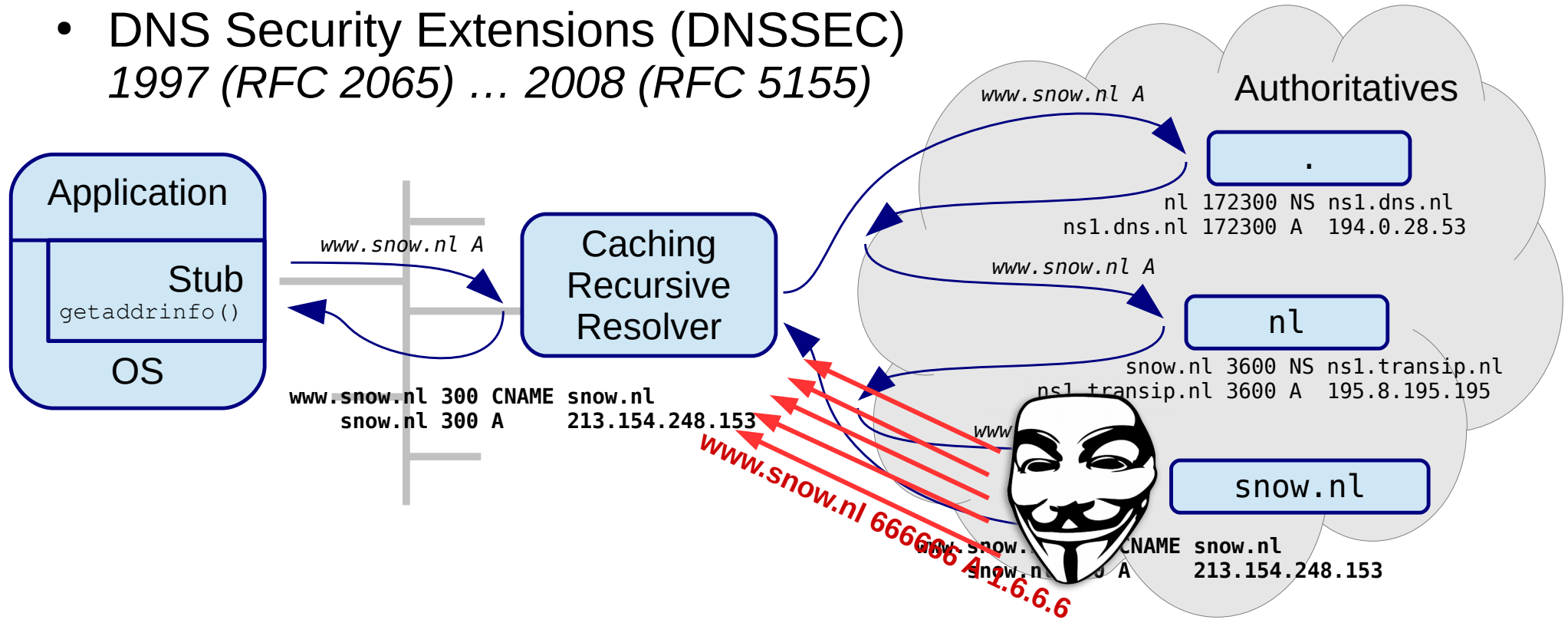
Domain Name System - schaal

- **UDP** = Geen State op authoritatives
- **Caching** Recursive Resolvers:
 - Verminderen load naar authoritatives
 - Verminderen latency naar stub



Domain Name System - security

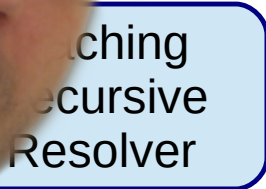
- Random bits (65.536 query ID * source ports) & **Caching** als security mechanisme
- DNS Security Extensions (DNSSEC)
1997 (RFC 2065) ... 2008 (RFC 5155)



Domain Name System

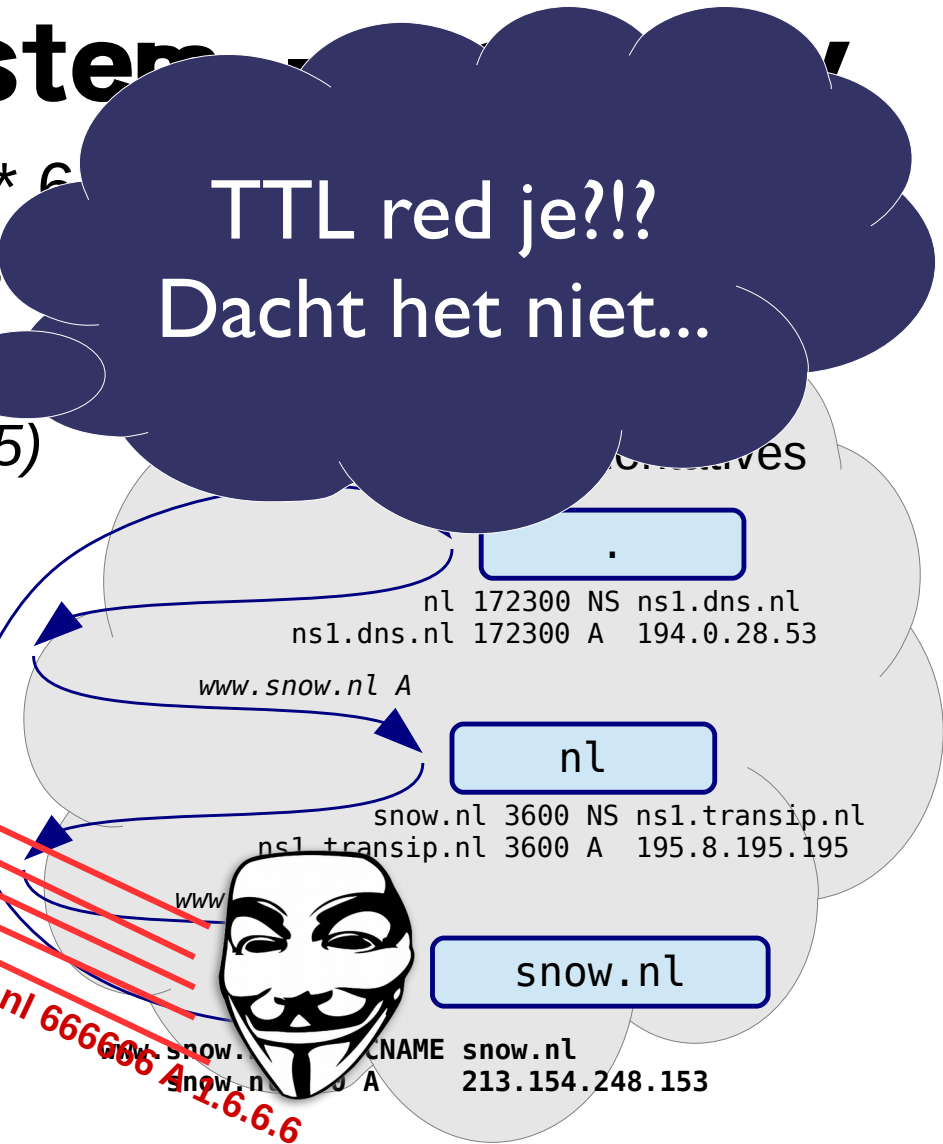
- Random bits (65.536 query ID * 65.536)
- **Caching** als security mechanism
- DNS Security Extensions (DNSSEC) RFC 5155

TTL red je?!?
Dacht het niet...

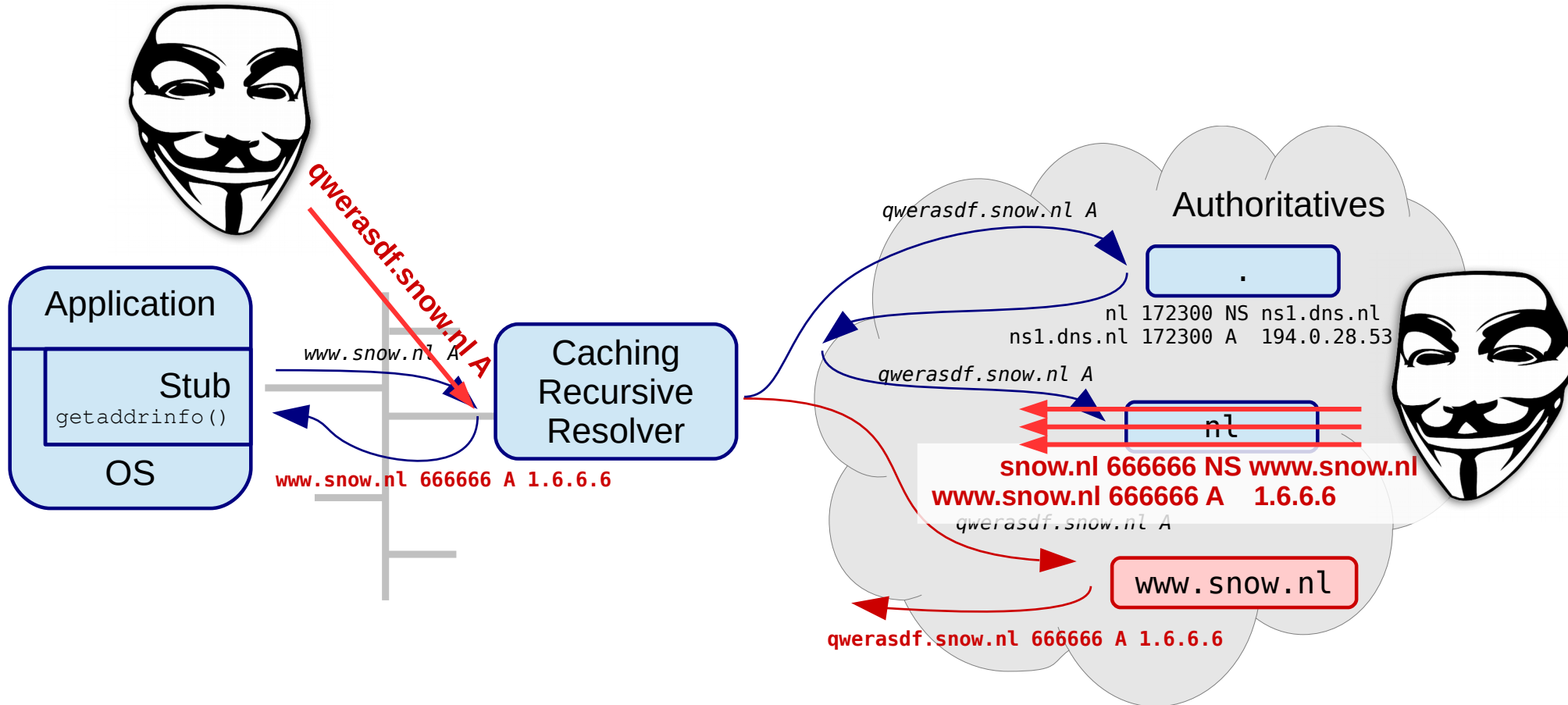


ME snow.nl
123.154.248.153

Security Popstar



Domain Name System - security



Domain Name System - security

# Bits	50% kans	5% kans	Methode
16	10 seconden	1 seconde	Query ID
26	2,8 uur	17 minuten	1024 source poorten
34	28 dagen	2,8 dagen	Alle source poorten + 2 bits server selectie
44	288444 dagen	2844,4 dagen	0x20 hack

Domain Name System – security

- Hulp bij spoofen van DNS antwoorden

Fragmentation Considered Poisonous

Amir Herzberg[†] and Haya Shulman[‡]

Dept. of Computer Science, Bar Ilan University

[†]amir.herzberg@gmail.com, [‡]haya.shulman@gmail.com

Abstract

Recent practical *poisoning* and *name-server block-*
ing attacks on standard DNS resolvers, by *off-path*,
adversaries. Our attacks exploit large DNS
responses that cause IP fragmentation; such long re-
sponses are increasingly common, mainly due to the use

in scenarios, where DNSSEC is partially or

sary that is able to send spoofed packets (but not to inter-
cept, modify or block packets). The most well known
is Kaminsky's DNS poisoning attack [21], which was
exceedingly effective against many resolvers at the time
(2008). Kaminsky's attack, and most other known DNS
poisoning attacks, allows the attacker to cause resolvers
to provide incorrect (poisoned) responses to DNS queries
of the clients, and thereby 'hijack' a domain name. We
refer to this type of attacks as *Domain hijacking DNS poi-*

Security
Rockstar

Domain Name System - security

- Hulp bij spoofen van DNS antwoorden

attacker ICMP frag needed → authoritative

Offsets	Octet	0				1				2				3																			
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	v4				IHL = 20				TOS				Total Length = 56				IP Header															
4	32	IPID								x DF MF				Frag Offset																			
8	64	TTL				Protocol = 1				IP Header Checksum																							
12	96	Source IP = 6.6.6.6																															
16	128	Destination IP = 2.2.2.2								ICMP Header																							
20	160	Type = 3				Code = 4					ICMP Checksum																						
24	192	Unused									MTU = 100																						
28	224	v4				IHL = 20				TOS				Total Length = 76				IP Header															
32	256	IPID								x DF MF				Frag Offset																			
36	288	TTL				Protocol = 17				IP Header Checksum																							
40	320	Source IP = 2.2.2.2								UDP Header																							
44	352	Destination IP = 7.7.7.7																															
48	384	Source Port = 53				Destination Port = 12345																											
52	416	Length = 56								UDP Checksum = 0																							

Security Rockstar

ent practical p
es on stand
adversaries.
hat cause

increasingly common, mainly due to the use

scenarios, where DNSSEC is partially or

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Domain Name System - security

- Hulp bij spoofen van DNS antwoorden

1^e fragment
authoritative → resolver

Offsets	Octet	0				1				2				3																			
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	v4		IHL = 20		TOS		Total Length = 85																									
4	32	IPID = 23456								x DF MF		Frag Offset = 0																					
8	64	TTL				Protocol = 17				IP Header Checksum																							
12	96	Source IP = 2.2.2.2																															
16	128	Destination IP = 7.7.7.7																															
20	160	Source Port = 53								Destination Port = 12345																							
24	192	Length = 65								UDP Checksum = 0x14de																							
28	224	TXID = 76543				QR		Opcode = 0		AA		TC		RD		RA		Z		RCODE = 0													
32	256	Question Count = 1								Answer Record Count = 1																							
36	288	Authority Record Count = 0								Additional Record Count = 1																							
40	320	4				m				a				i																			
44	352	l				4				v				i																			
48	384	c				t				2				i																			
52	416	m				0				Type = A																							
56	448	Class = IN								Name (Pointer)																							
60	480	Type = A								Class = IN																							
64	512	TTL																															

2^e fragment
attacker → resolver

Offsets	Octet	0				1				2				3																			
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	v4		IHL = 20		TOS		Total Length = 85																									
4	32	IPID = 23456								x DF MF		Frag Offset = 48																					
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12	96	Source IP = 2.2.2.2																															
16	128	Destination IP = 7.7.7.7																															
20	160	Data Length = 4								IPv4 Address																							
24	192	= 2.2.2.2								Name = 0				Type																			
28	224	= OPT				UDP Payload Size = 4096								EXTENDED-RCODE = 0																			
32	256	Version = 0				DO		Z				Data Length																					
36	288	= 0																															

server block-
by off-path,
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uch long re-
due to the use

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scenarios, where DNSSEC is partially or

Domain Name System - security

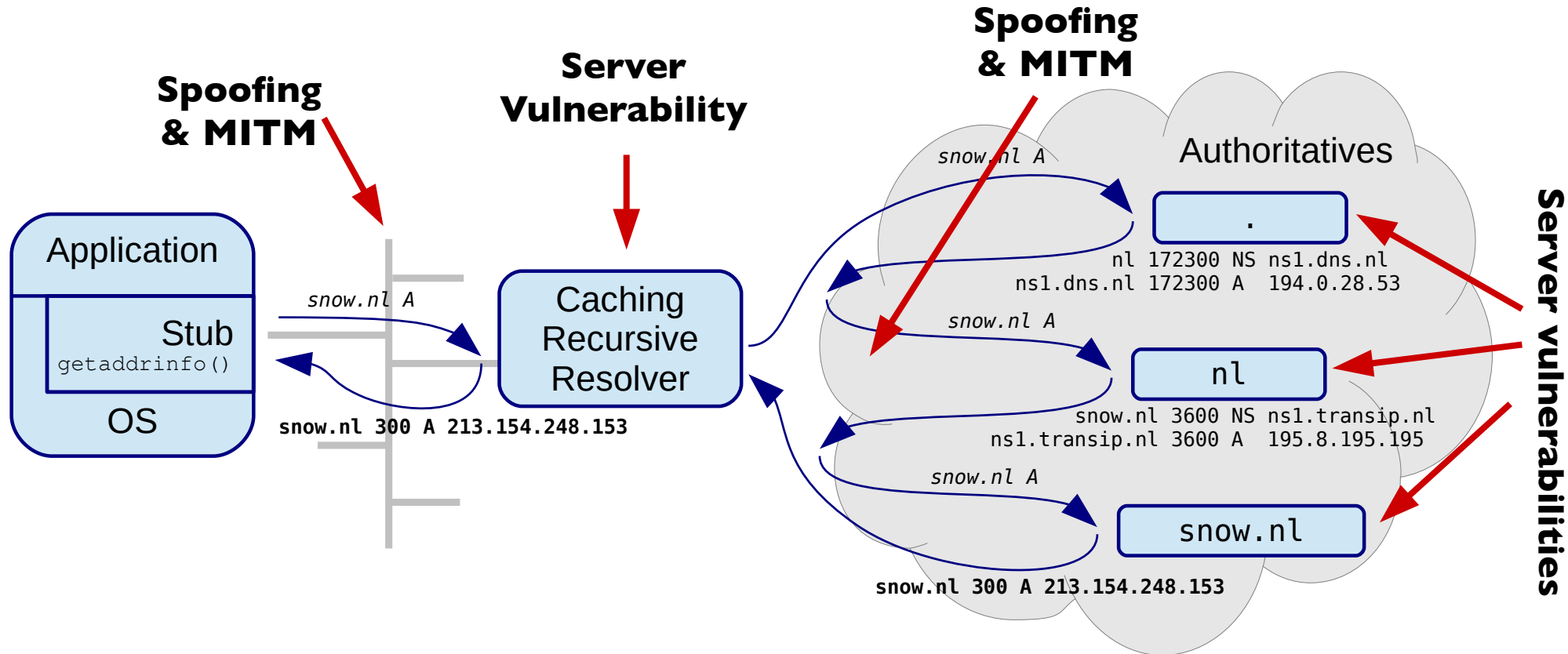
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Domain Name System - security

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5	0 seconden	0 seconden	IP ID
69	2.928.370.544 jaar	2.928.370.544 jaar	IPv6 /64 source adres

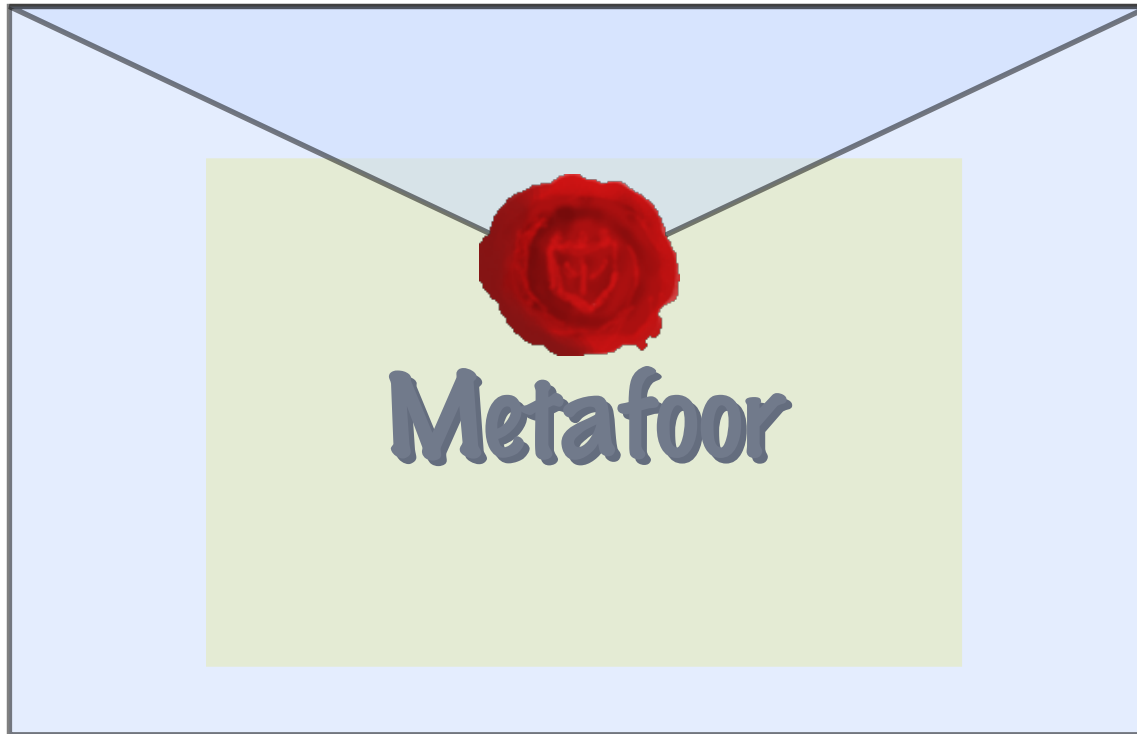
Domain Name System - security

- 't is niet alleen spoofing



DNS Security Extensions (DNSSEC)

- end-to-end security bovenop DNS

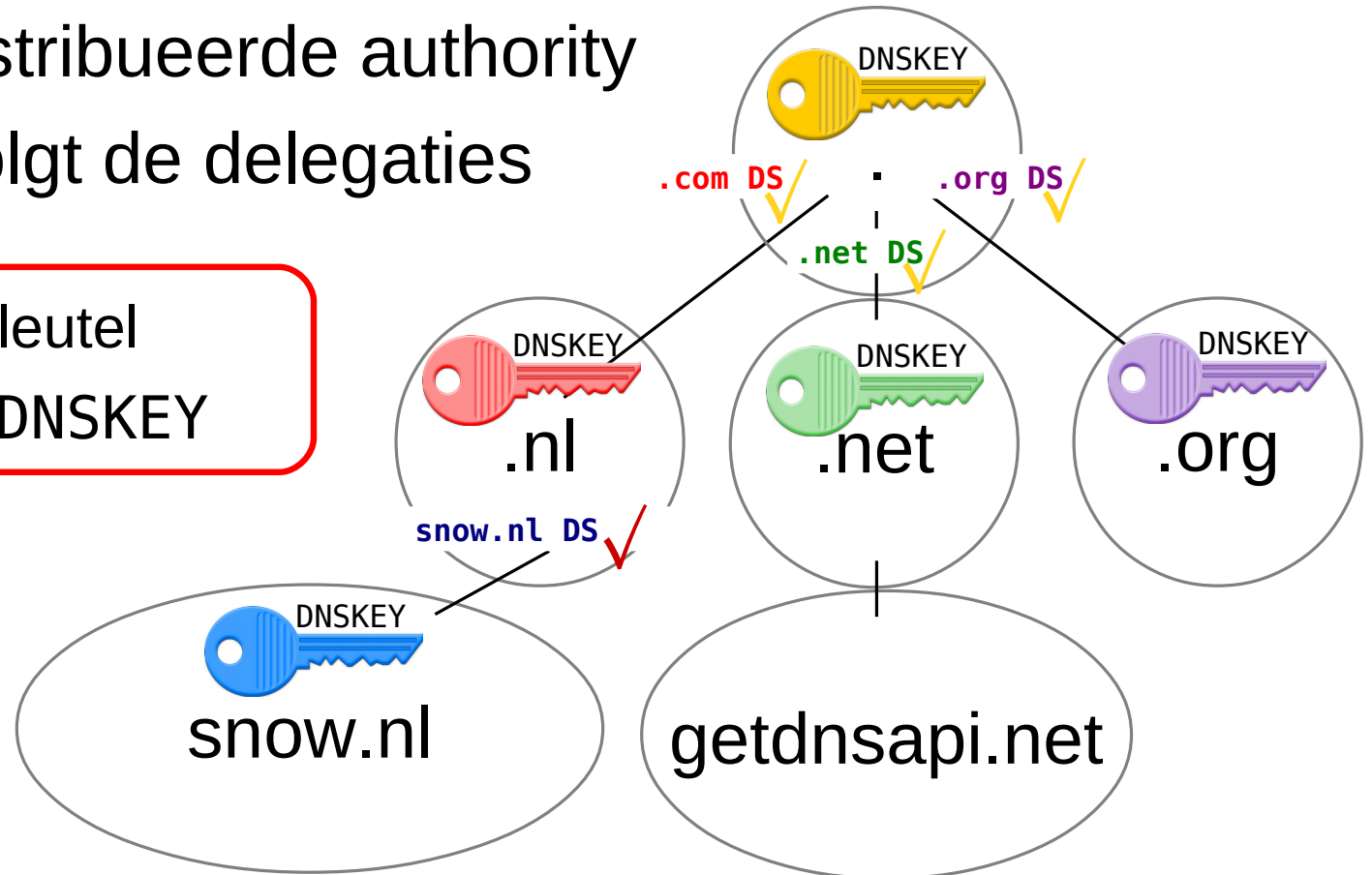


DNS Security Extensions (DNSSEC)

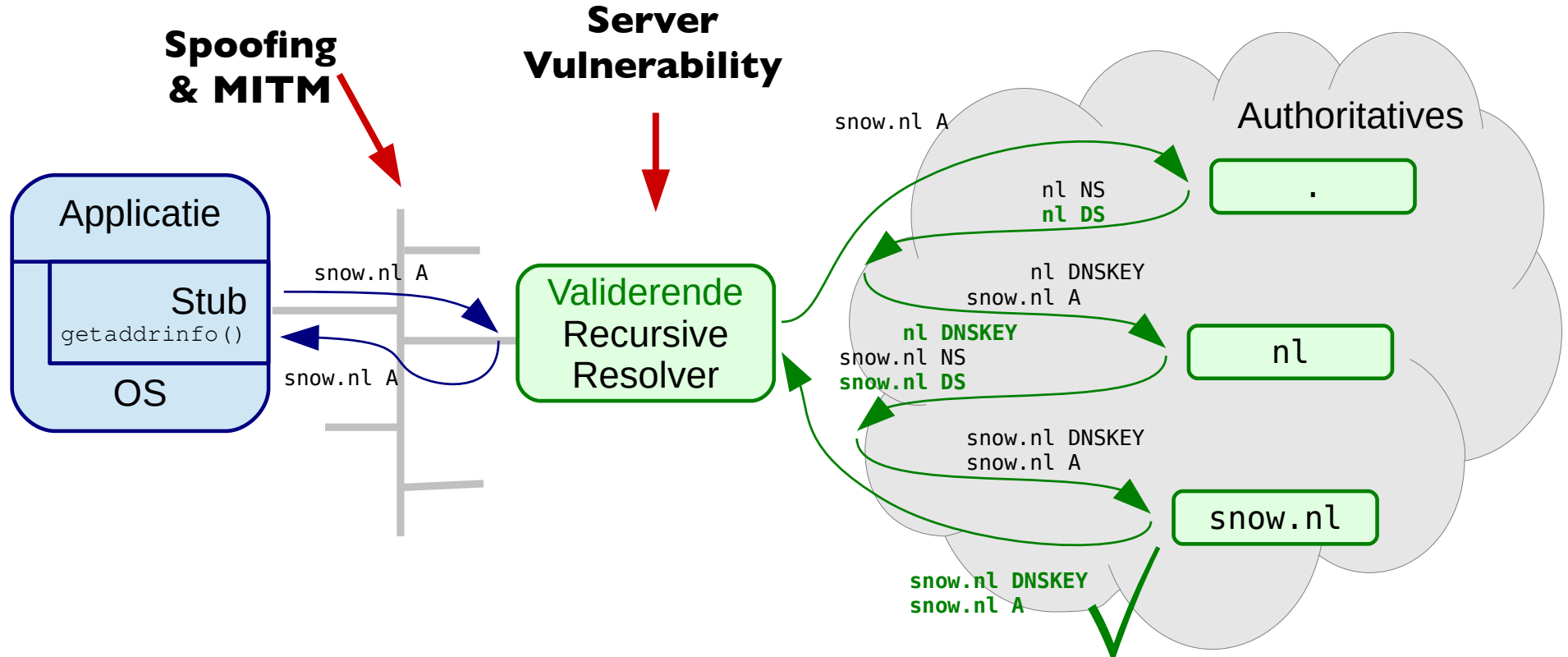
Chain of trust

- Zones met gedistribueerde authority
- Chain of trust volgt de delegaties

- DNSKEY Publieke sleutel
- DS Hash van DNSKEY

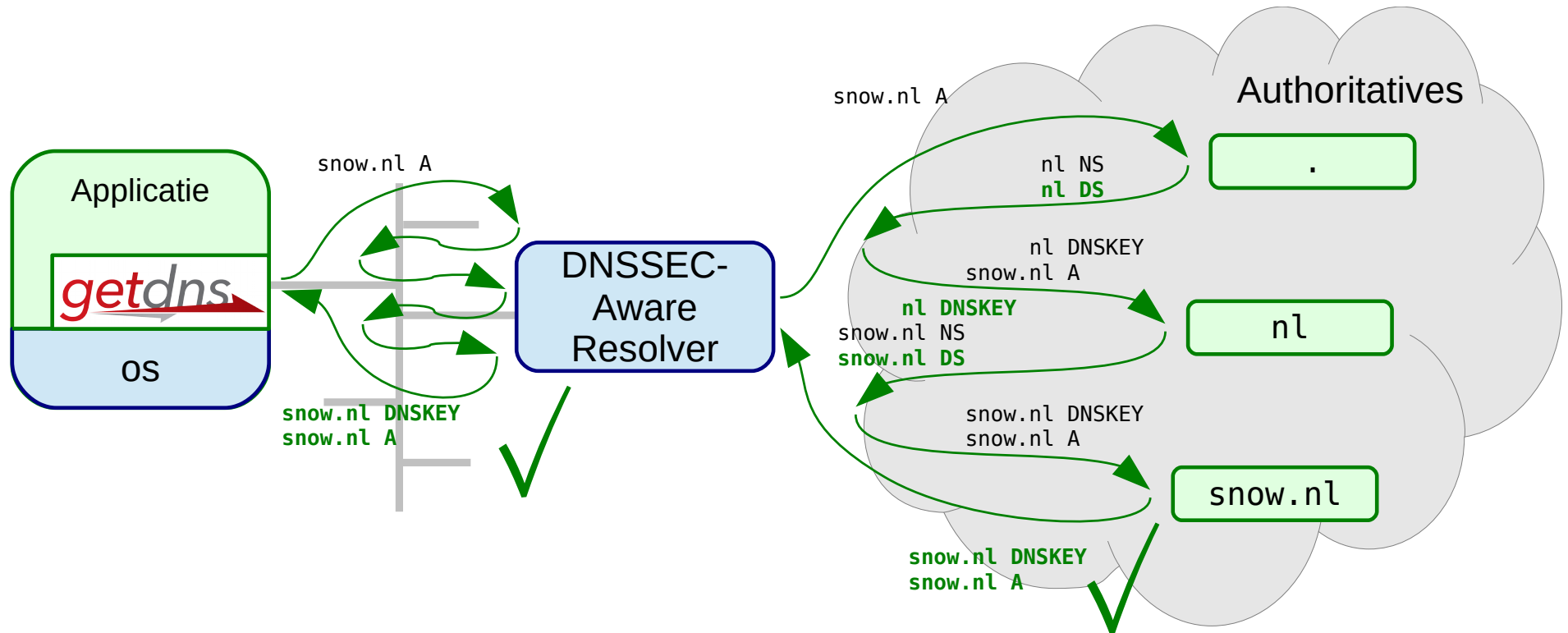


DNS Security Extensions (DNSSEC) Validatie



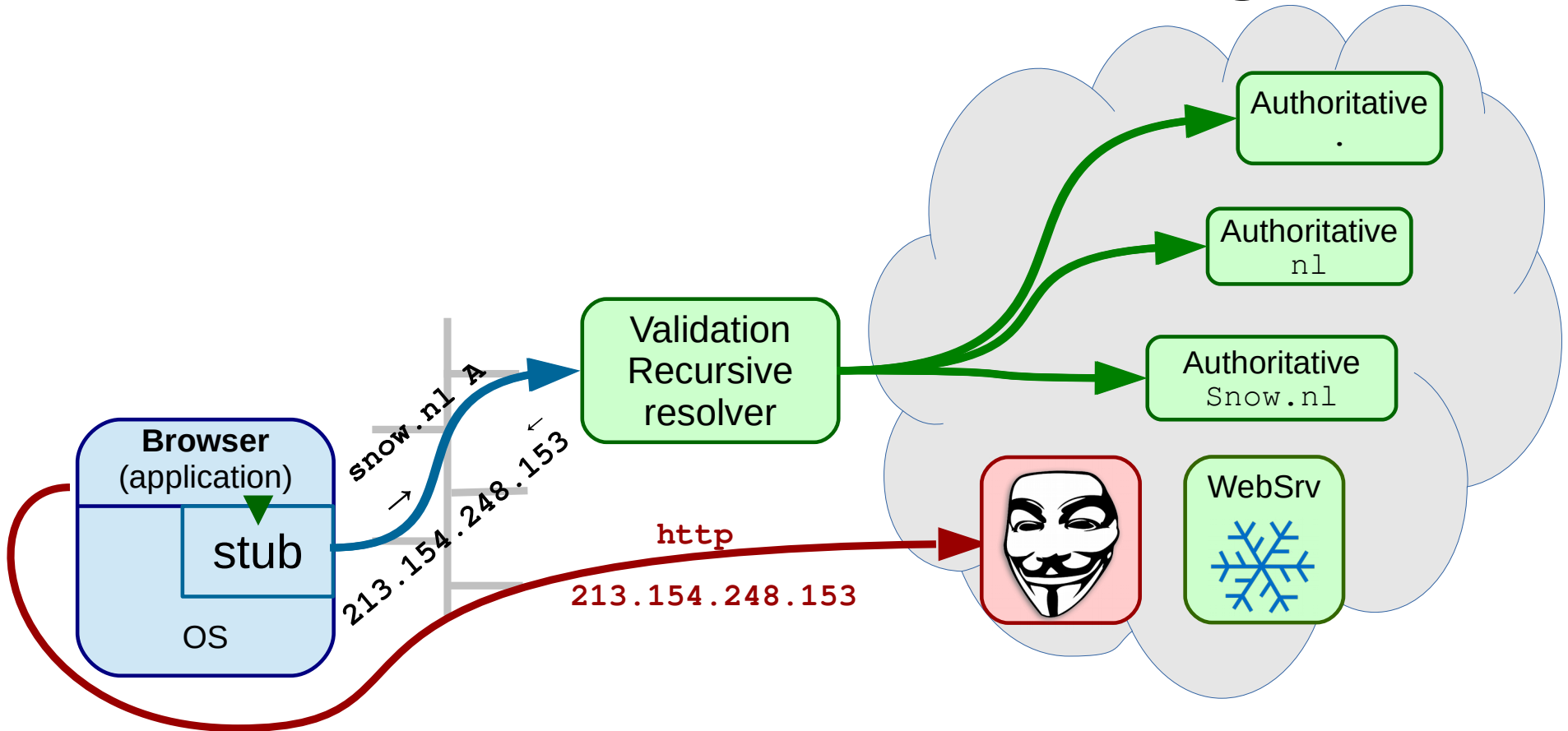
DNS Security Extensions (DNSSEC)

end-to-end validatie



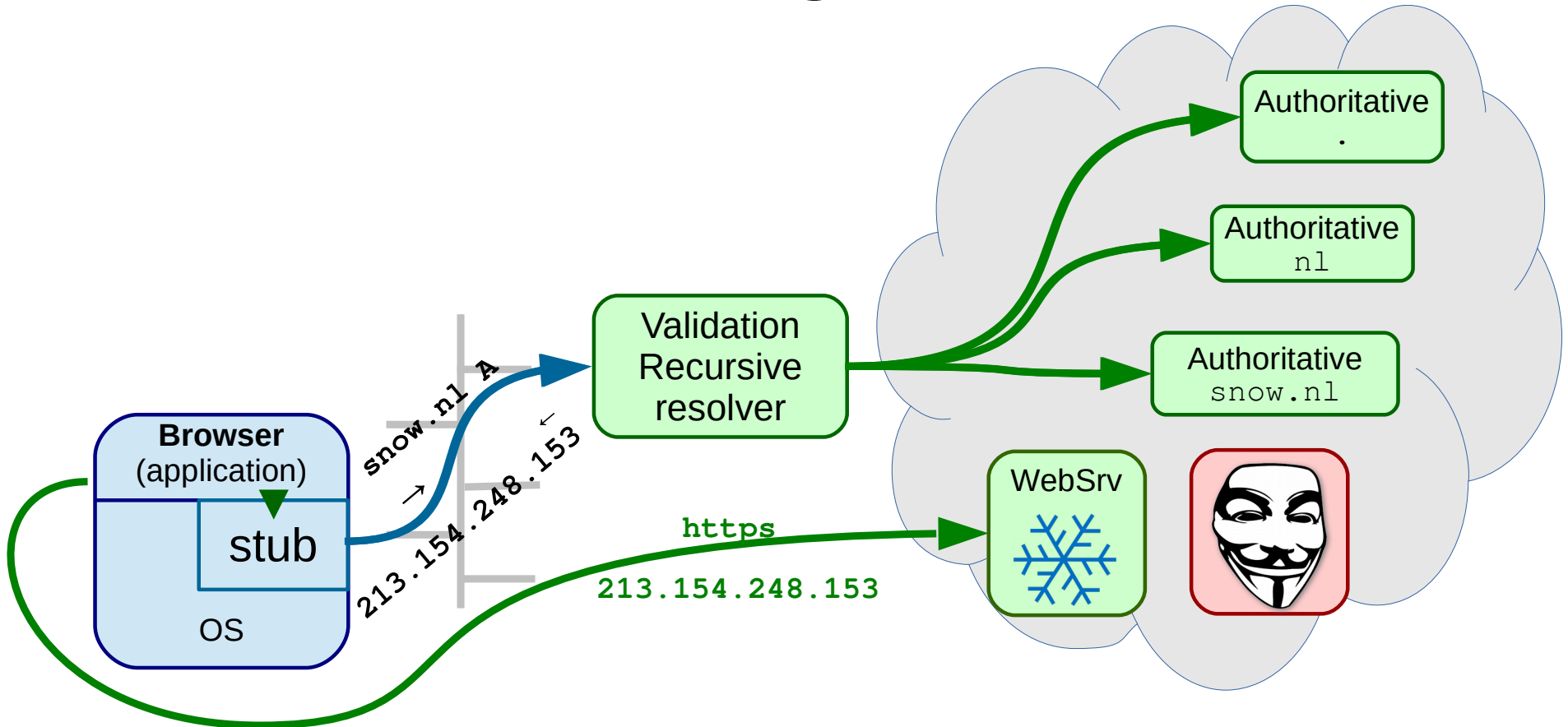
DNS Security Extensions (DNSSEC)

beschermt niet tegen MITM



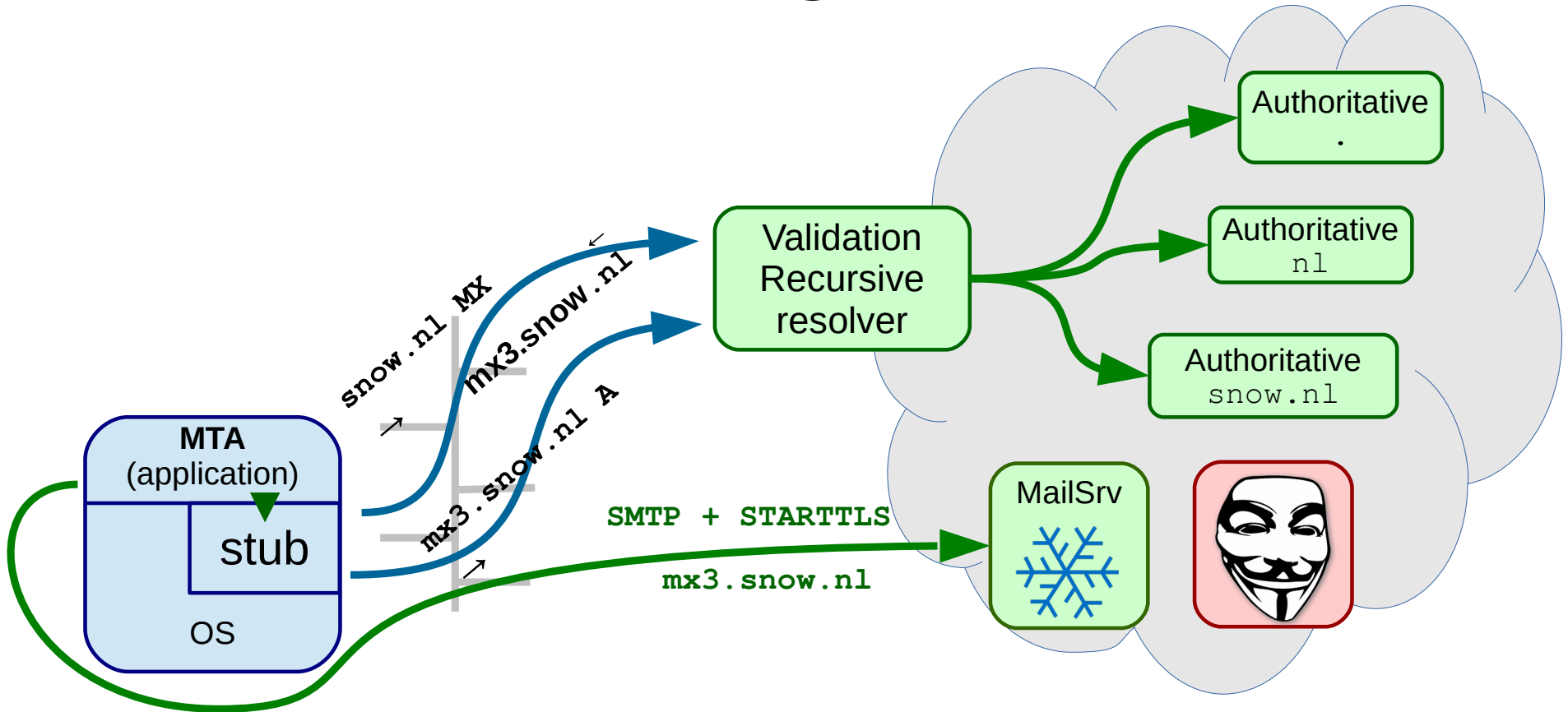
DNS Security Extensions (DNSSEC)

beschermt niet tegen MITM – TLS wel!



DNS Security Extensions (DNSSEC)

toch nodig voor DNS referrals



DNSSEC voor Applicaties

voor TLS

- Transport Layer Security (TLS) gebruikt zowel asymmetrische als symmetrische cryptografie
- Een symmetrische sleutel wordt versleuteld verstuurd samen met de publieke sleutel van de andere kant

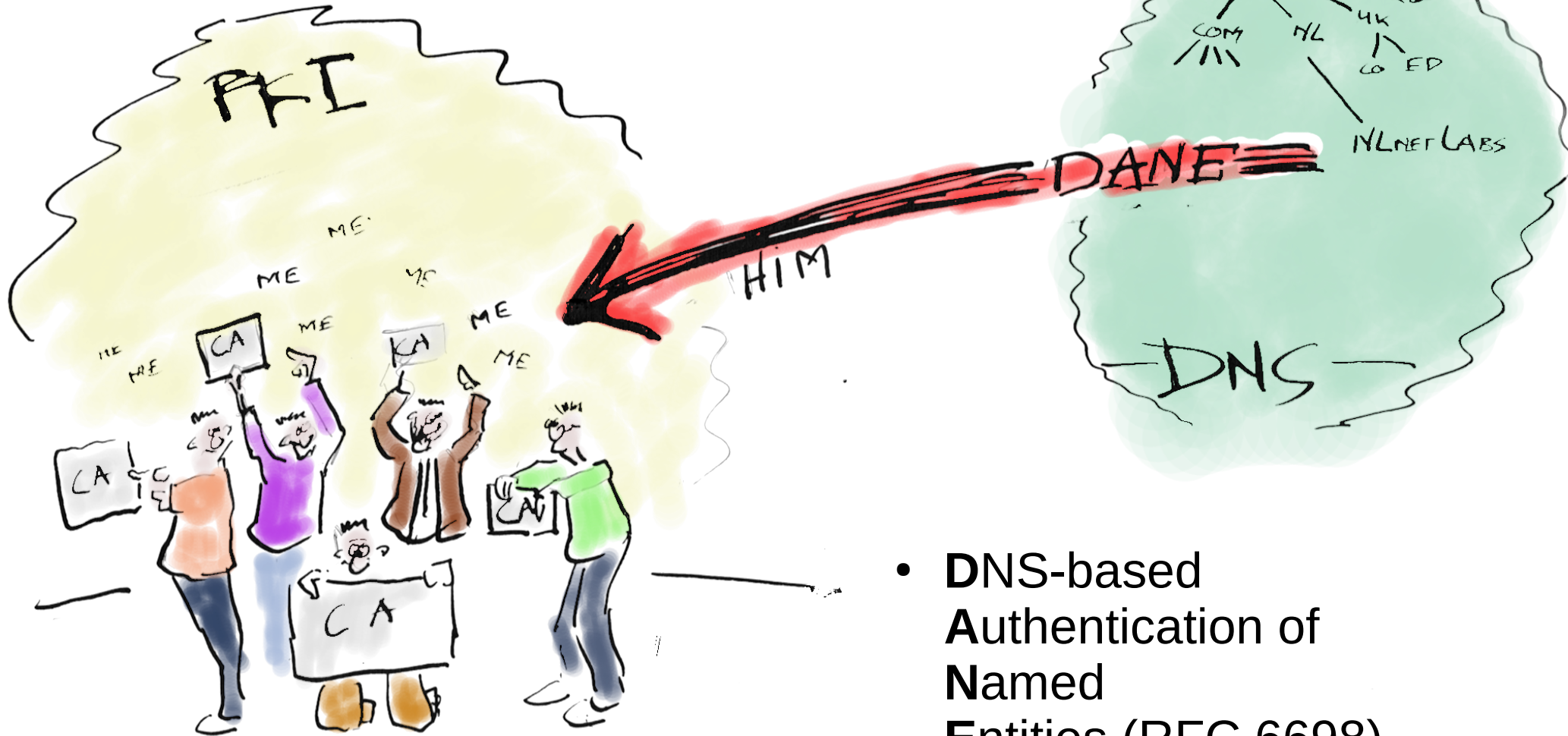
- Hoe wordt die publieke sleutel geverifieerd?

TLS zonder DNSSEC



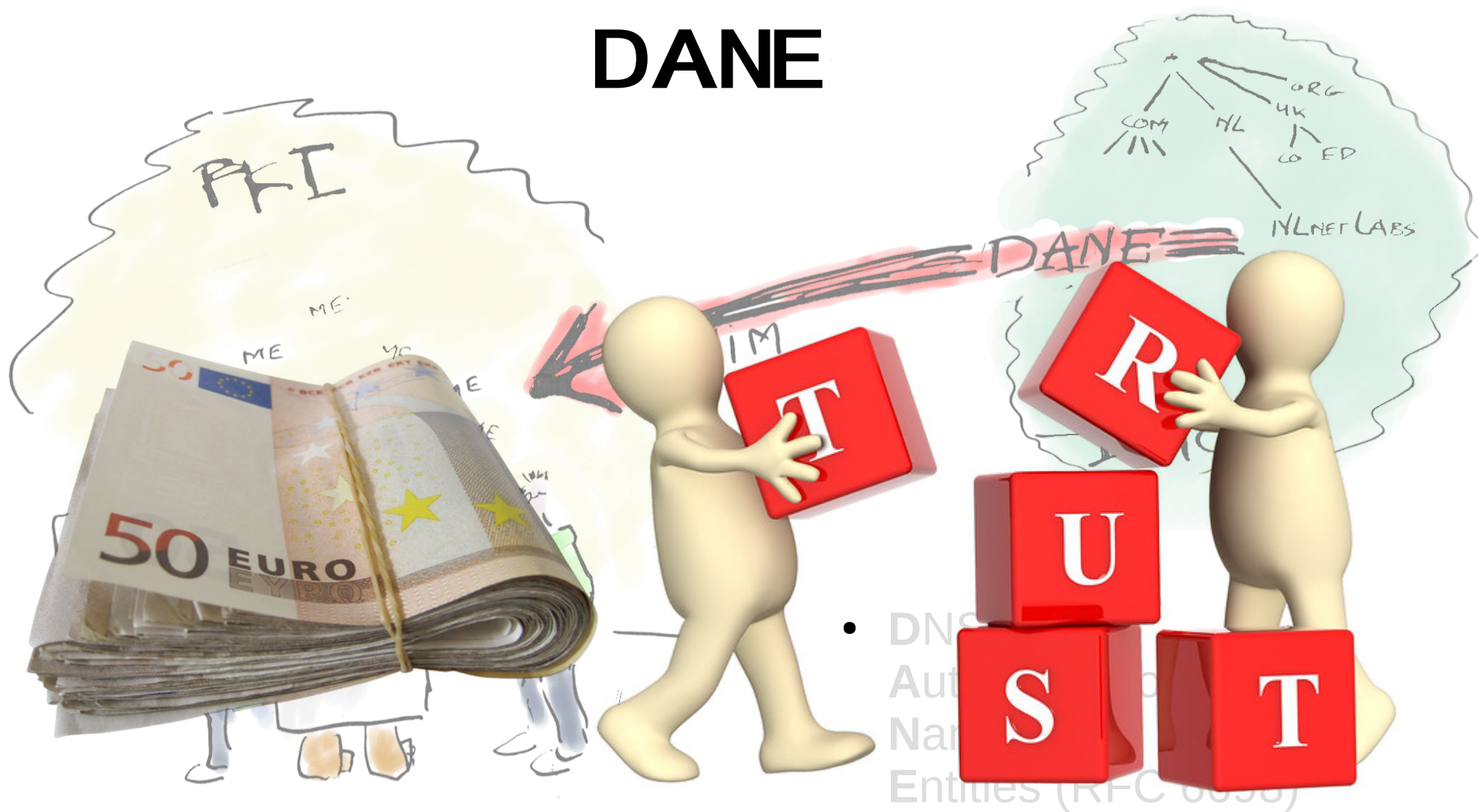
- Door de Certificate Authorities in OS en/of browser
- Elke CA is gemachtigd in te staan voor **elke** domein naam
- Er zijn meer dan 1500 CAs
(in 2010, zie <https://www.eff.org/observatory>)

DANE



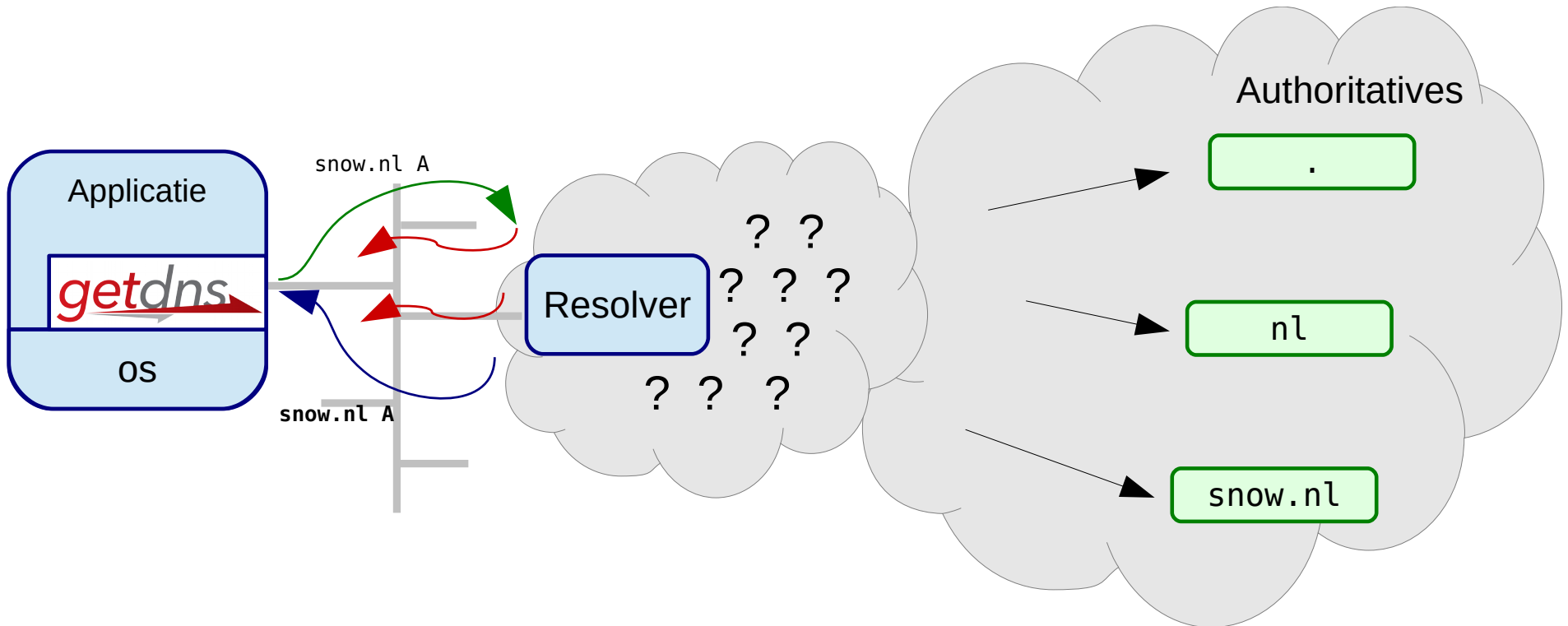
- DNS-based Authentication of Named Entities (RFC 6698)

DANE



DNS Security Extensions (DNSSEC)

end-to-end validatie in de praktijk



DNS Security Extensions (DNSSEC)

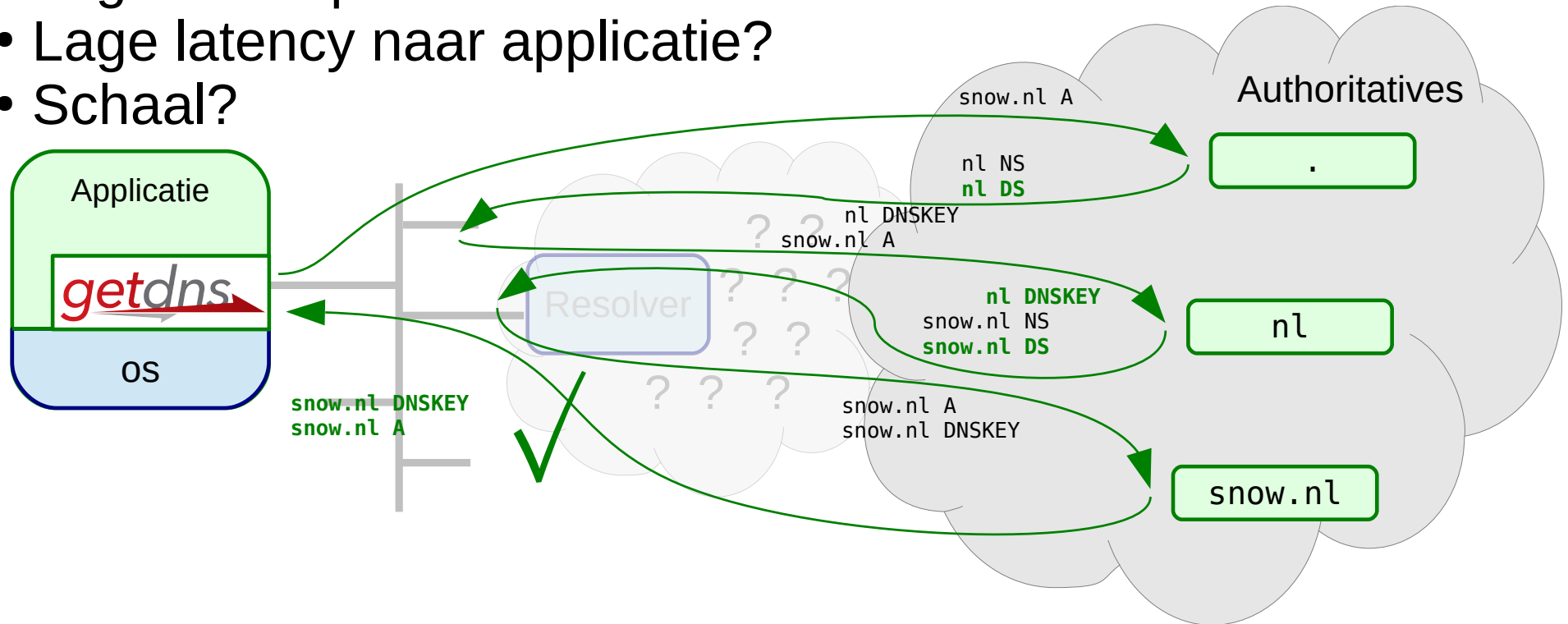
end-to-end validatie in de praktijk



DNS Security Extensions (DNSSEC)

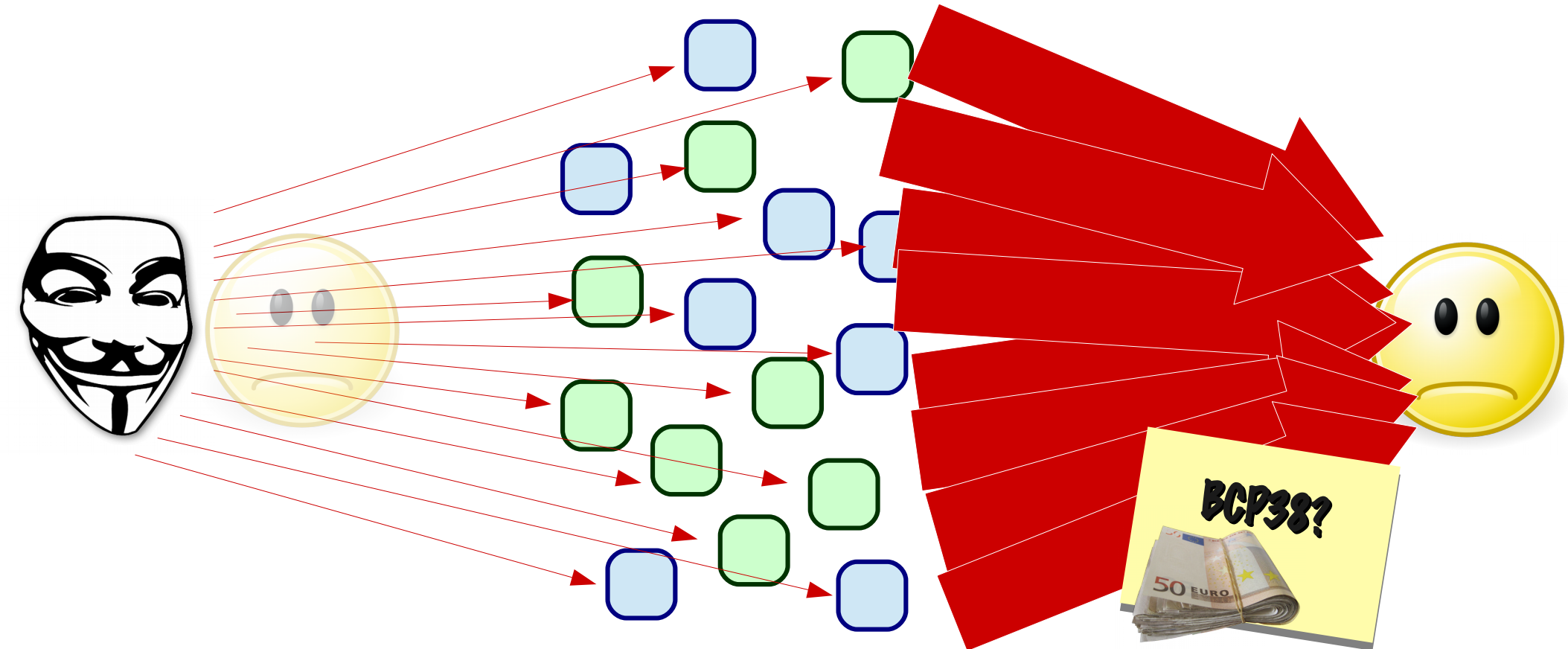
end-to-end validatie in de praktijk

- Lage load op de authoritaties?
- Lage latency naar applicatie?
- Schaal?



DNS Security Extensions (DNSSEC)

consequentie van UDP erger met DNSSEC



Privacy

maart 2011: I-D

Privacy Considerations
for Internet Protocols

juni 2013: ██████████ Revelations
[Morecowbell](#)

juli 2013: RFC6973
Privacy Considerations
for Internet Protocols

mei 2014: RFC7258
Pervasive Monitoring
is an Attack



**Privacy
Folk Singer**

Overall Encryptie

ations
protocols

juni 2013: **Revelations**

[Morecowbell](#)

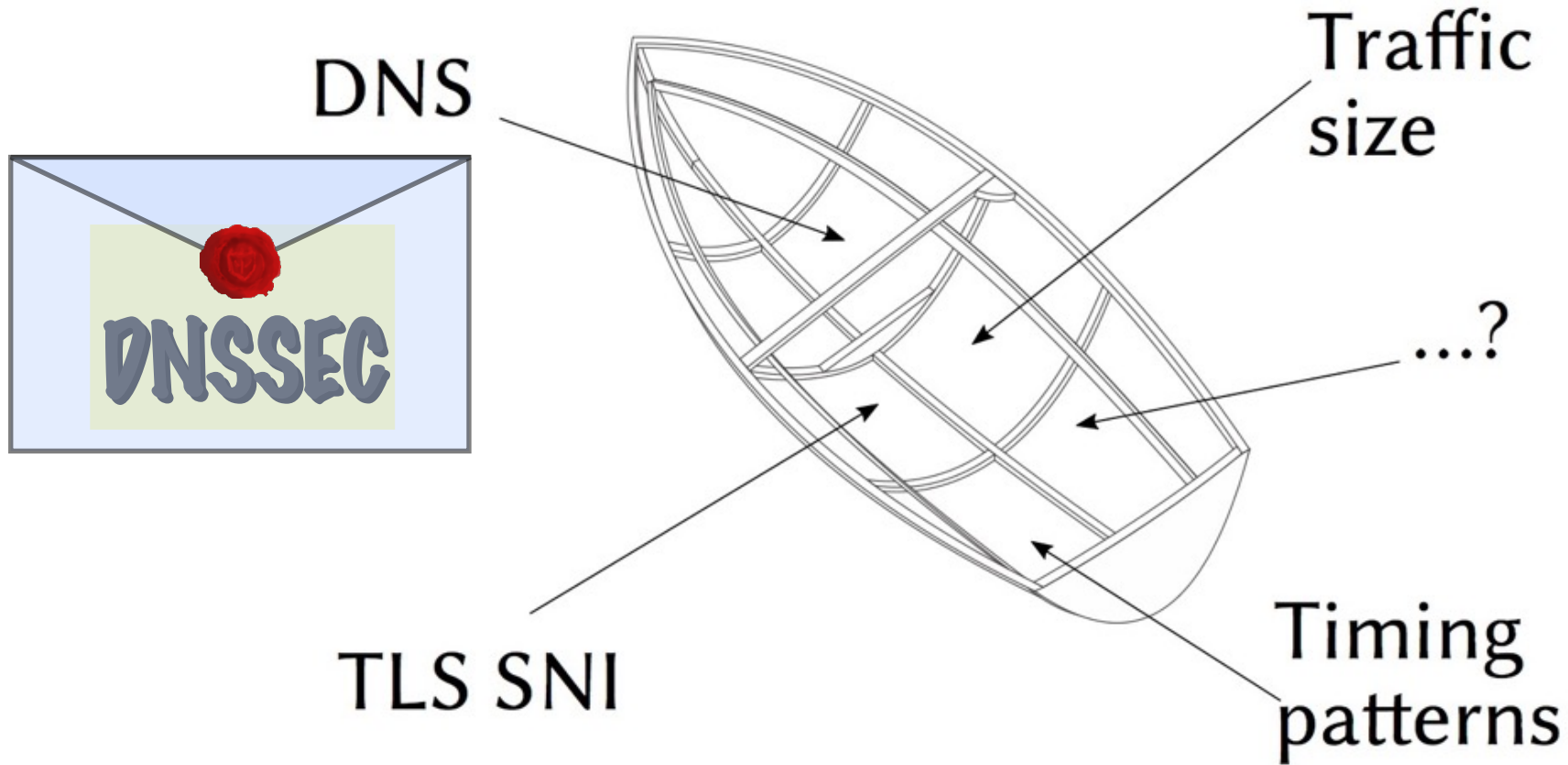
juli 2013 : RFC6973
Privacy Considerations
for Internet Protocols

mei 2014: **RFC7258**
**Pervasive Monitoring
is an Attack**



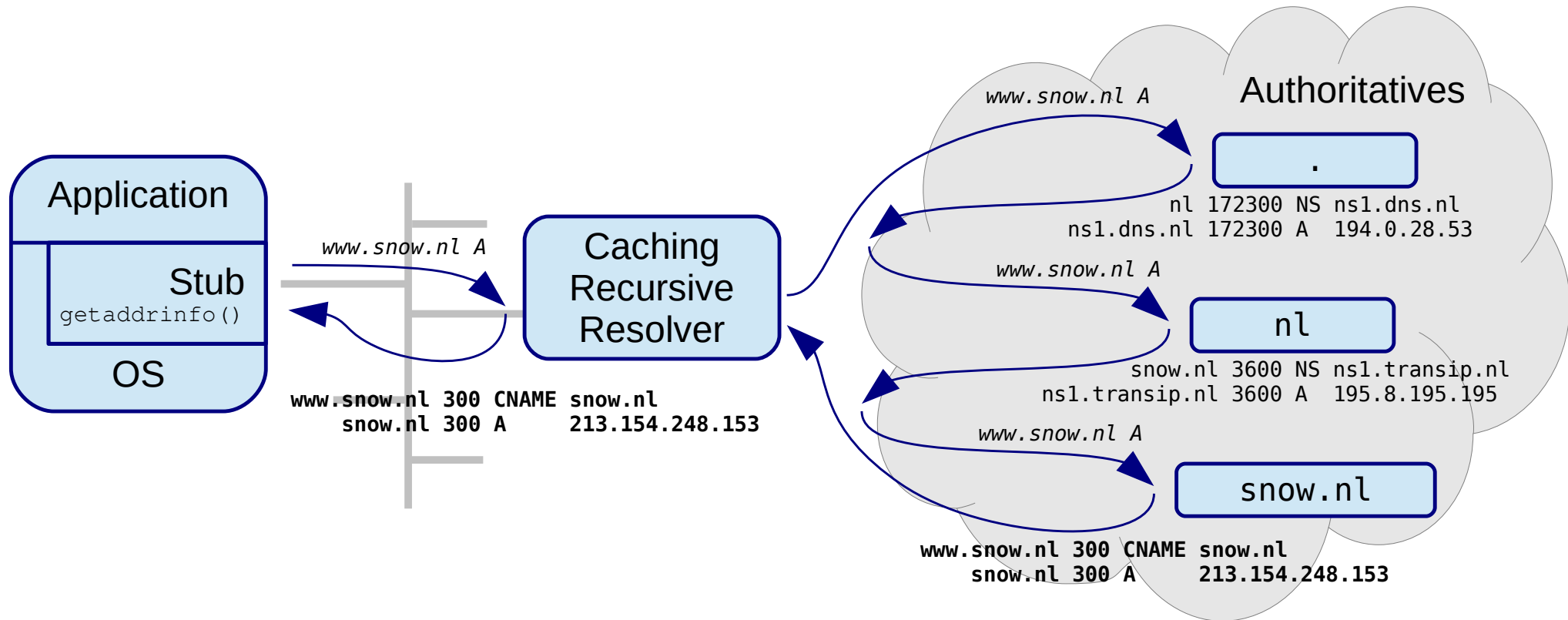
**Privacy
Folk Singer**

Privacy



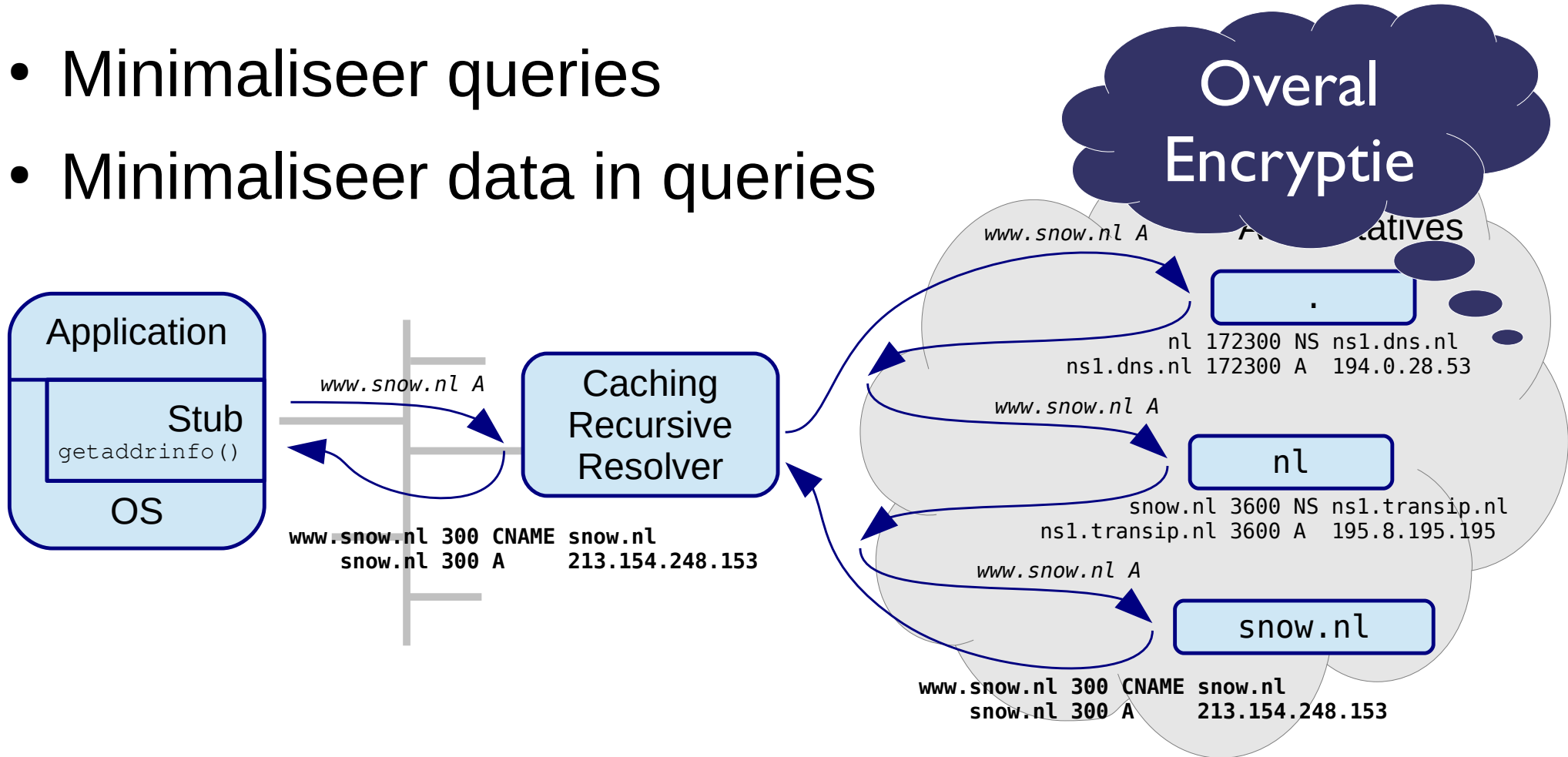
- NSA's [Morecowbell](#) op DNS gebaseerde monitoring systeem

Privacy issues met DNS



Privacy issues met DNS

- Minimaliseer queries
- Minimaliseer data in queries



Privacy issues met DNS

minimaliseer queries – local root

- RFC 7706 -
Running a Root Server
Local to a Resolver

```
auth-zone:  
  name: "."  
  master: 199.9.14.201  
  master: 192.33.4.12  
  master: 199.7.91.13  
  master: 192.5.5.241  
  master: 192.112.36.4  
  master: 193.0.14.129  
  master: 192.0.47.132  
  master: 192.0.32.132  
  fallback-enabled: yes  
  for-downstream: no  
  for-upstream: yes
```

```
"unbound.conf"
```



unbound

Privacy issues met DNS

minimaliseer queries – local auth zone

- RFC 7706 -
Running a Root Server
Local to a Resolver
- Kan ook voor andere
authoritative servers

```
auth-zone:  
  name: "se"  
  master: zonedata.iis.se  
  zonefile: "se.zone"  
  fallback-enabled: yes  
  for-downstream: no
```

```
"unbound.conf"
```



unbound

Privacy issues met DNS

minimaliseer queries – aggressive NSEC

- RFC8198 -
Aggressive NSEC

```
$ dig @k.root-servers.net snow. +norec +dnssec

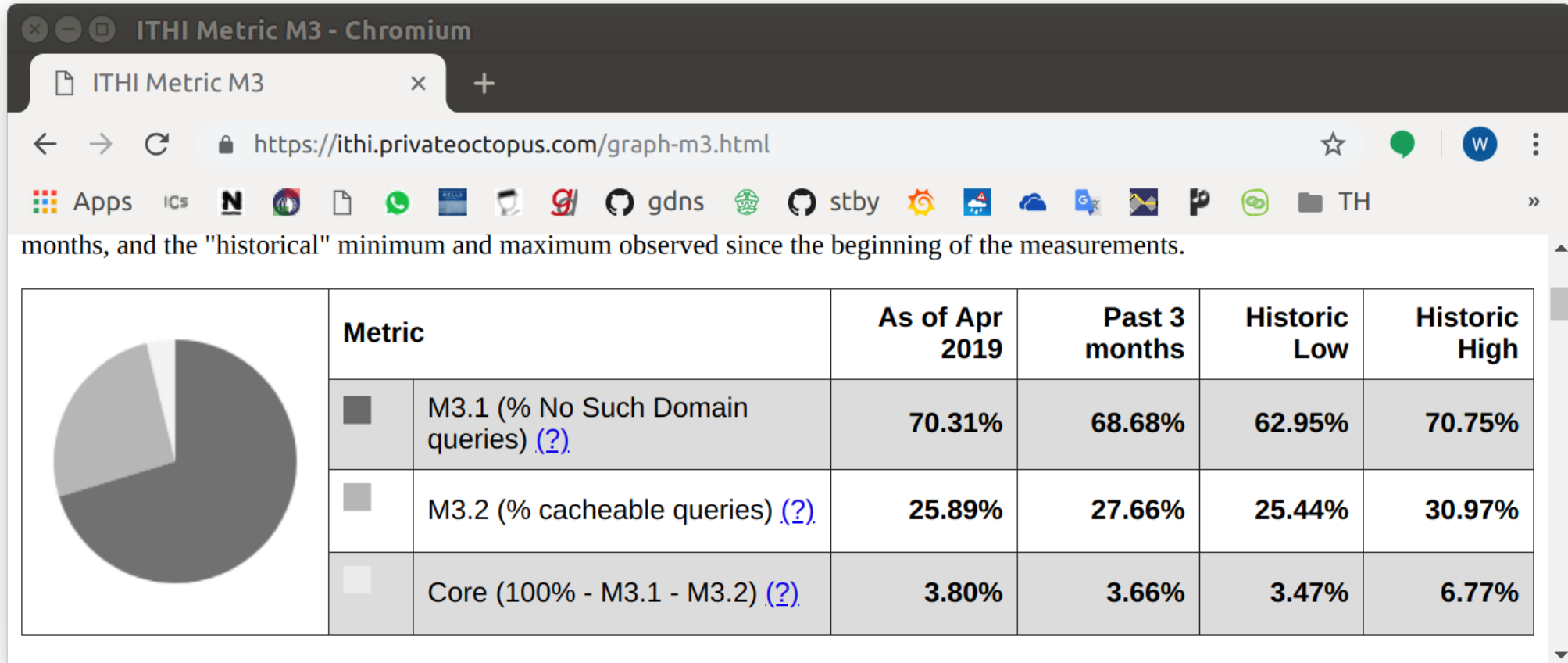
;; ->HEADER<<- opcode: QUERY, rcode: NXDOMAIN, id:
;; flags: qr aa ; QUERY: 1, ANSWER: 0, AUTHORITY: 6
;; QUESTION SECTION:
;; snow. IN  A

;; AUTHORITY SECTION:
sncf.      86400 IN NSEC so. NS DS RRSIG NSEC
sncf.      86400 IN RRSIG NSEC 8 1 86400 ...

.          86400 IN NSEC aaa. NS SOA RRSIG NSEC DNSKEY
.          86400 IN RRSIG NSEC 8 0 86400 ...

;; Query time: 2 msec
```

Privacy issues met DNS minimaliseer queries – aggressive NSEC



Privacy issues met DNS

minimaliseer queries – aggressive NSEC

- RFC8198 -
Aggressive NSEC

```
server:  
  aggressive-nsec: yes
```

```
"unbound.conf"
```



unbound

Privacy issues met DNS

minimaliseer queries – serve stale

- [draft-ietf-dnsop-serve-stale](#)
- Privacy aspect en/of Performance aspect

```
server:  
  serve-expired: yes  
  serve-expired-ttl: 300  
  serve-expired-ttl-reset: yes
```

```
"unbound.conf"
```

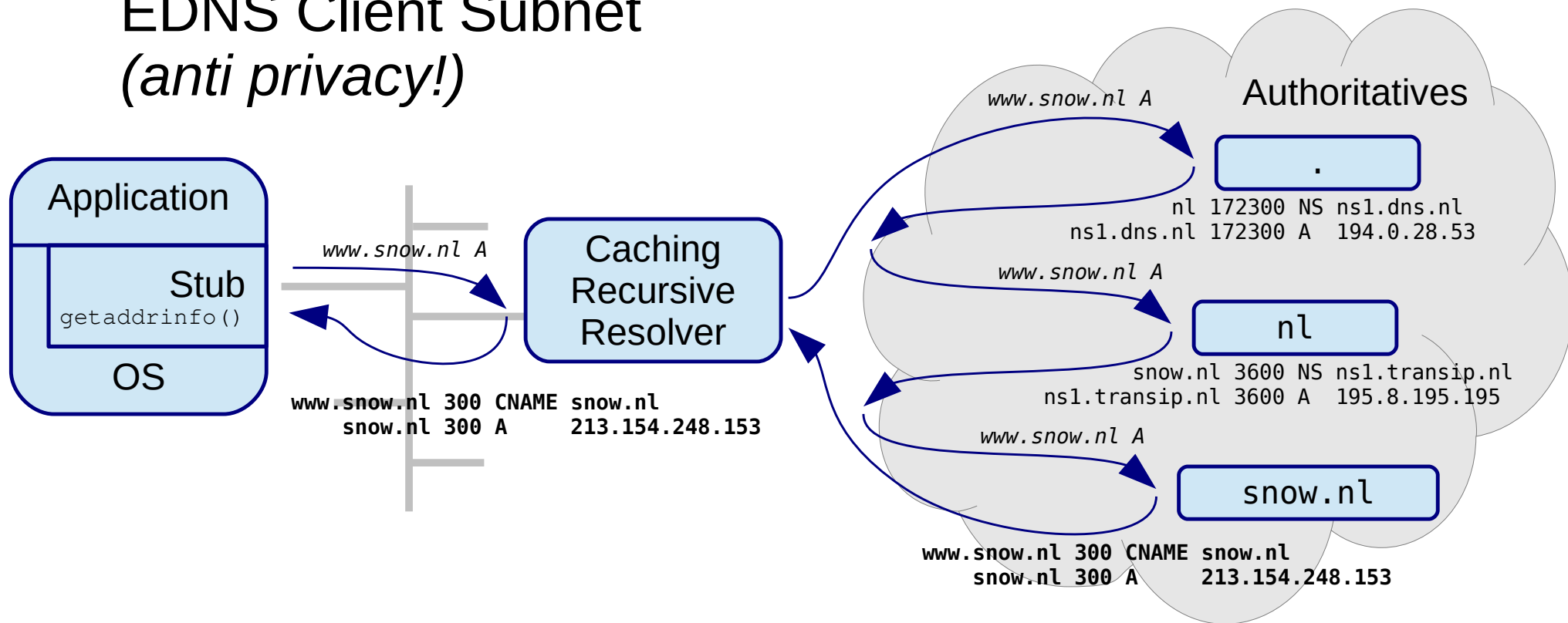


unbound

Privacy issues met DNS

minimaliseer data in queries – ECS

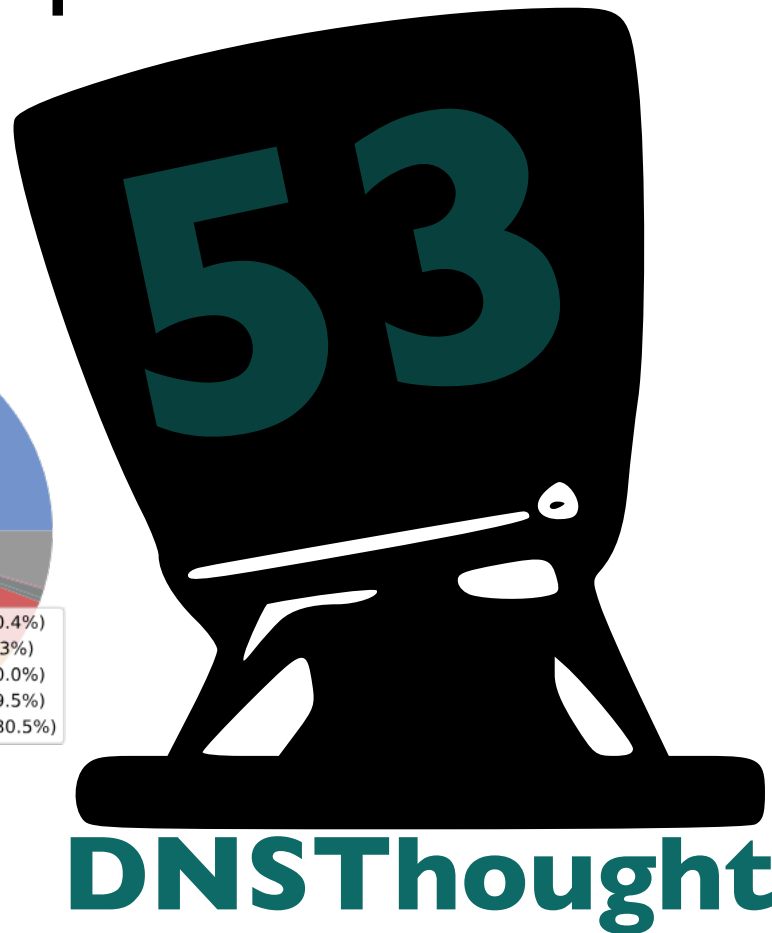
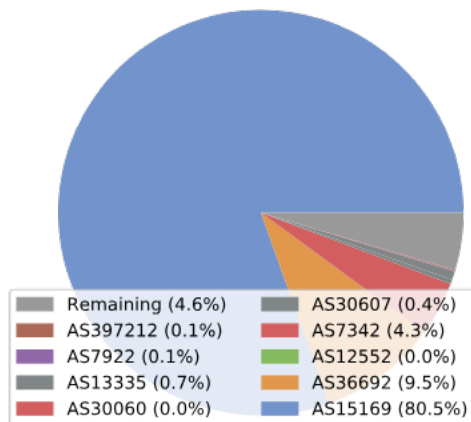
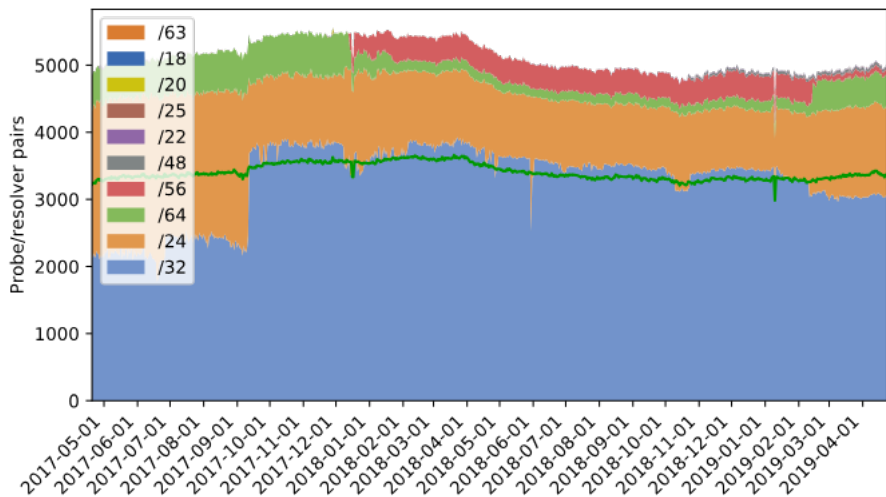
- RFC7871 -
EDNS Client Subnet
(*anti privacy!*)



Privacy issues met DNS

minimaliseer data in queries – ECS

- RFC7871 -
EDNS Client Subnet
(*anti privacy!*)



Privacy issues met DNS

minimaliseer data in queries – ECS priv.

- RFC7871 -
EDNS Client Subnet
sectie 7.1.2:
“ A SOURCE PREFIX-LENGTH value of 0 means that the Recursive Resolver MUST NOT add the client's address information to its queries. ”

 unbound respecteert dit

- Google respecteert dit

 OpenDNS respecteert dit niet

```
# EDNS0 option for ECS client privacy  
# as described in Section 7.1.2 of  
# https://tools.ietf.org/html/rfc7871
```

```
edns_client_subnet_private : 1
```

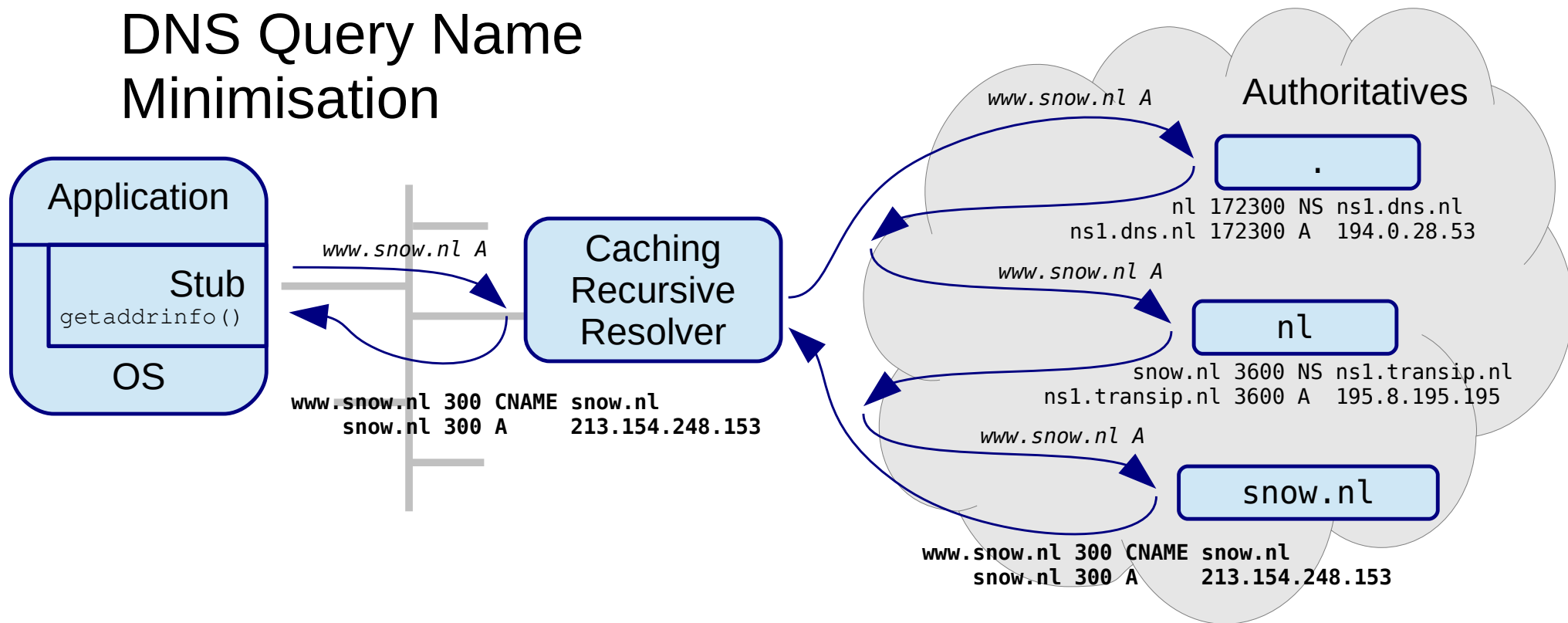
"stubby.yml"



Privacy issues met DNS

minimaliseer data in queries – qname min

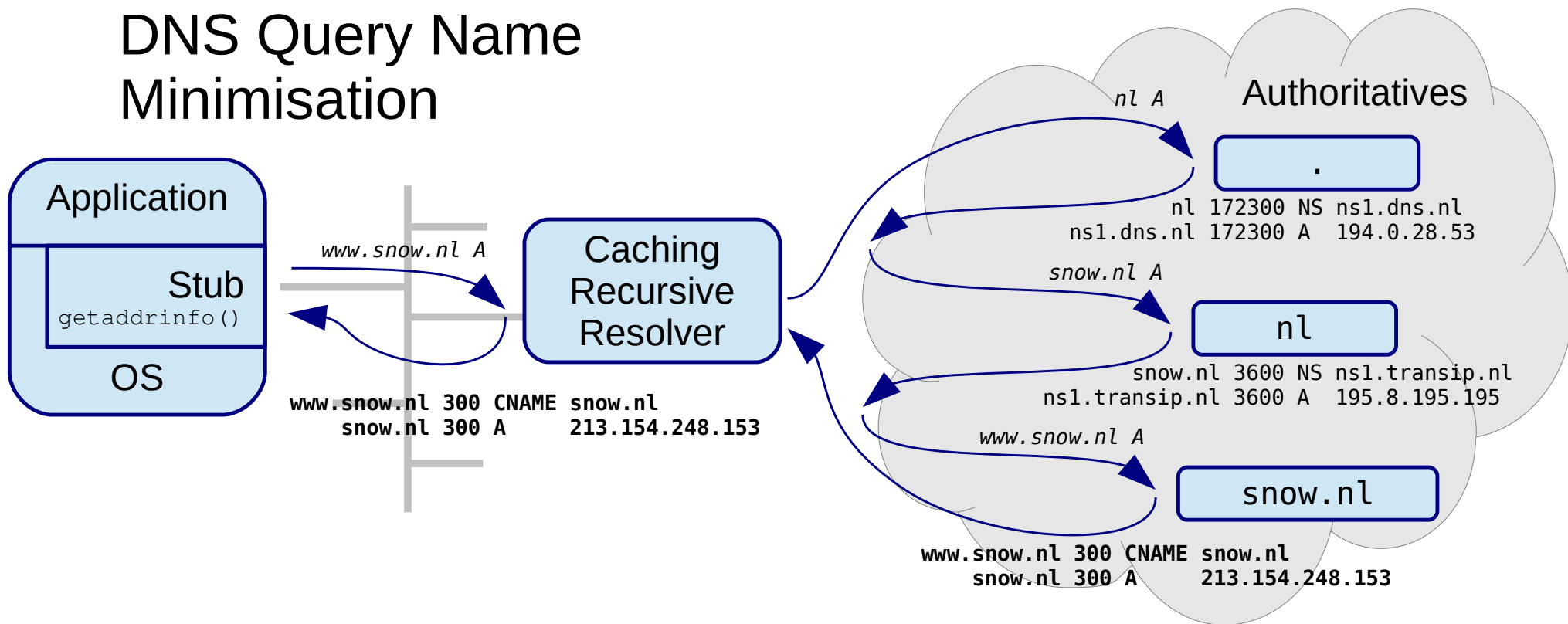
- Zonder RFC7816 -
DNS Query Name
Minimisation



Privacy issues met DNS

minimaliseer data in queries – qname min

- Met RFC7816 -
DNS Query Name
Minimisation



Privacy issues met DNS

minimaliseer data in queries – qname min

- RFC7816 -
DNS Query Name
Minimisation

```
server:  
  qname-minimisation: yes  
  qname-minimisation-strict: no
```

```
"unbound.conf"
```

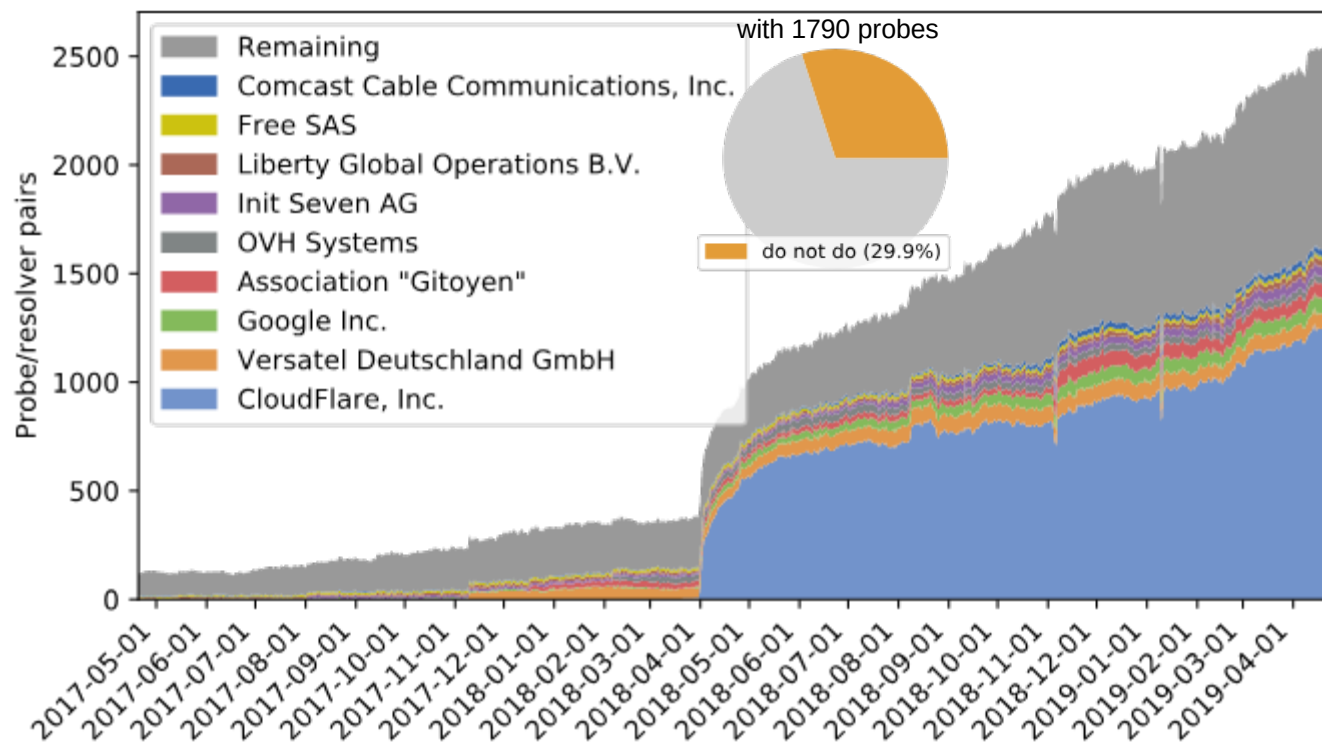


unbound

Privacy issues met DNS

minimaliseer data in queries – qname min

- RFC7816 - DNS Query Name Minimisation



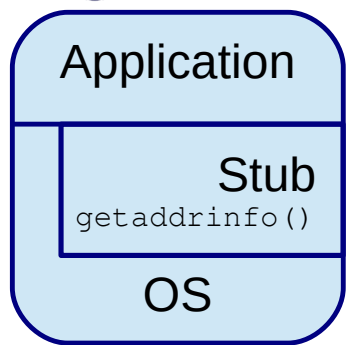
ITHI: 20.6% gemeten op de root

Privacy issues met DNS

Overall
Encryptie

minimaliseer (data in) queries

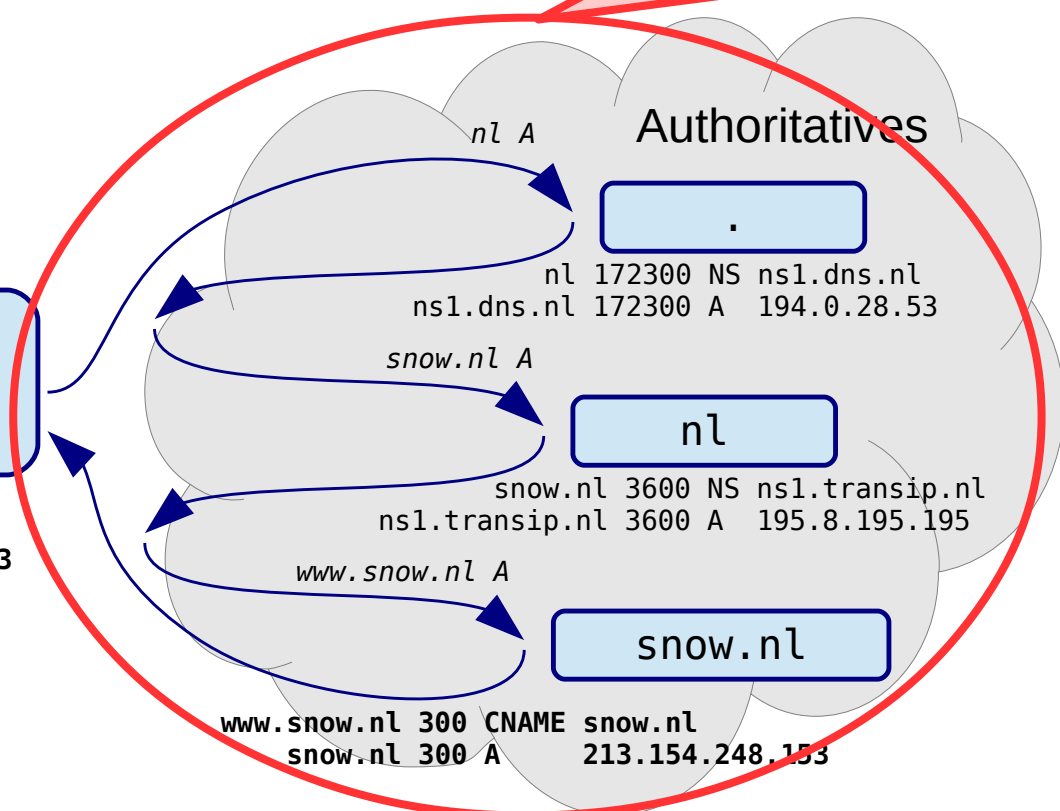
MITM, s
Eavesdroppers



`www.snow.nl A`

`www.snow.nl 300 CNAME snow.nl`
`snow.nl 300 A 213.154.248.153`

Caching
Recursive
Resolver

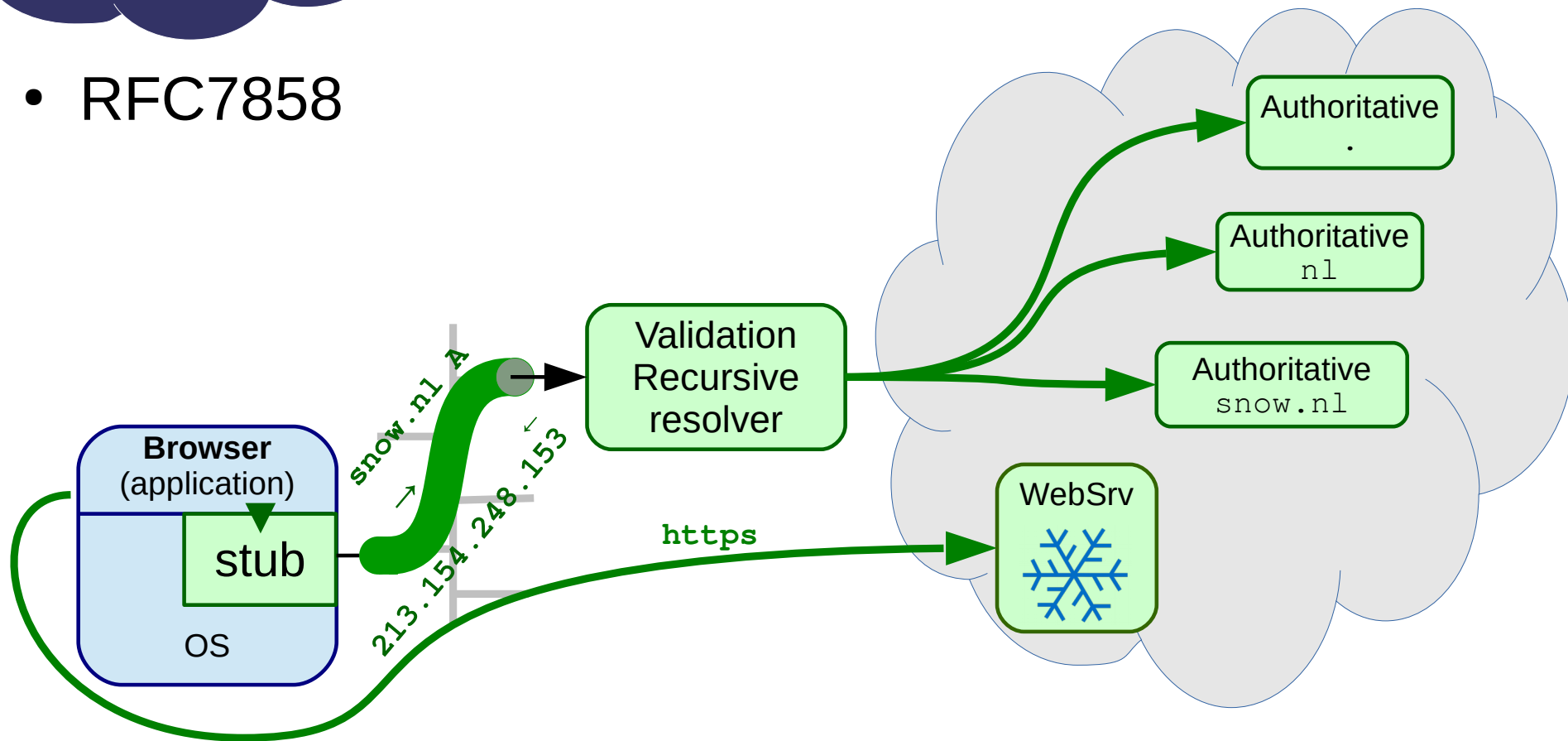


Overall
Encryption

Privacy issues met DNS

DNS over TLS (DoT)

- RFC7858

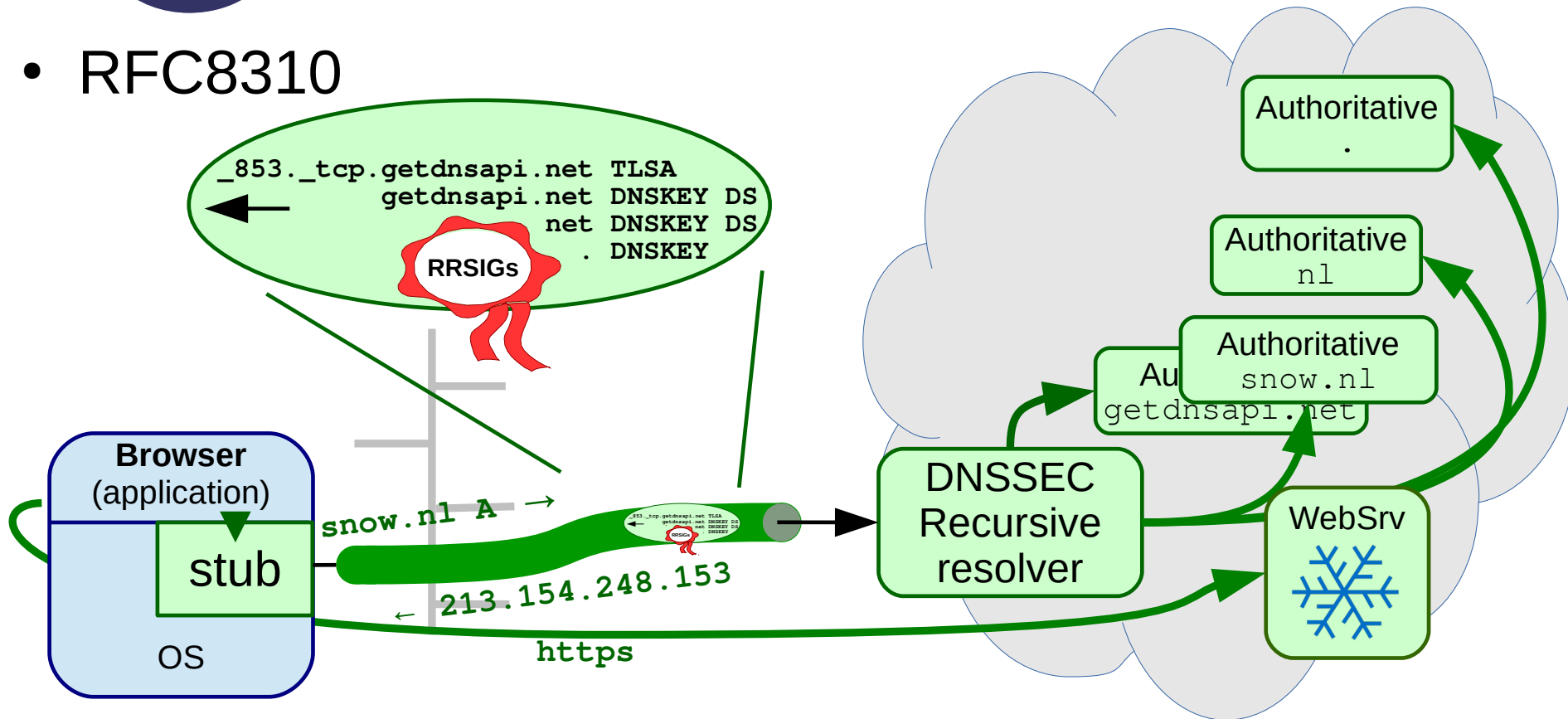


Overall
Encryptie

Privacy issues met DNS

DNS over TLS (DoT)

- RFC8310



Overall
Encryptie

Privacy issues met DNS

DNS over TLS (DoT)

```
server:  
  tls-service-key: "privkey.pem"  
  tls-service-pem: "fullchain.pem"  
  tls-port: 853
```

"unbound.conf"



```
round_robin_upstreams: 1  
  
upstream_recursive_servers:  
## Quad 9  
  - address_data: 9.9.9.9  
    tls_auth_name: "dns.quad9.net"  
## Cloudflare  
  - address_data: 1.1.1.1  
    tls_auth_name: "cloudflare-dns.com"  
## Google  
  - address_data: 8.8.8.8  
    tls_auth_name: "dns.google"
```

"stubby.yml"

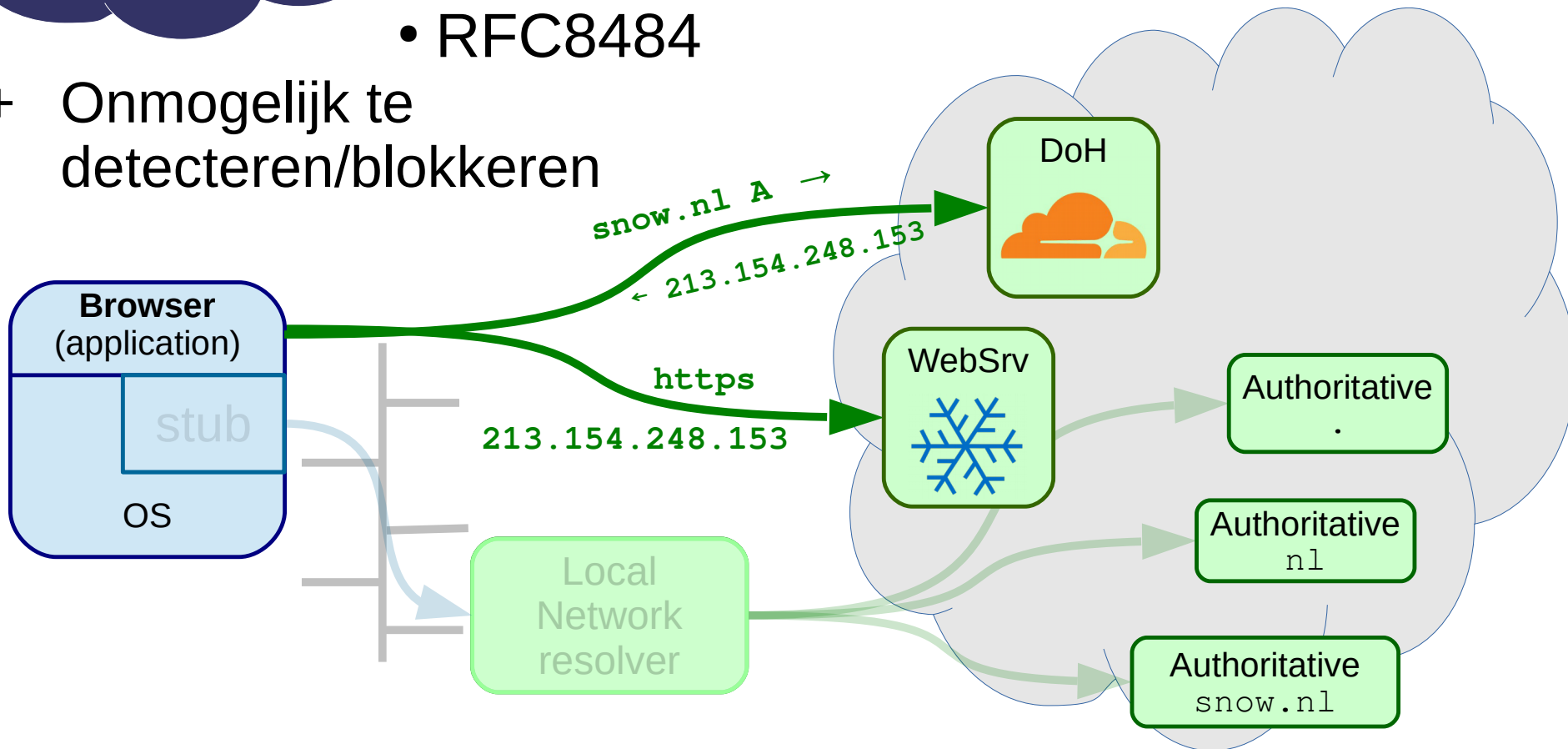


Overall
Encryptie

Privacy issues met DNS DNS over HTTPS (DoH)

- RFC8484

- + Onmogelijk te detecteren/blokkeren

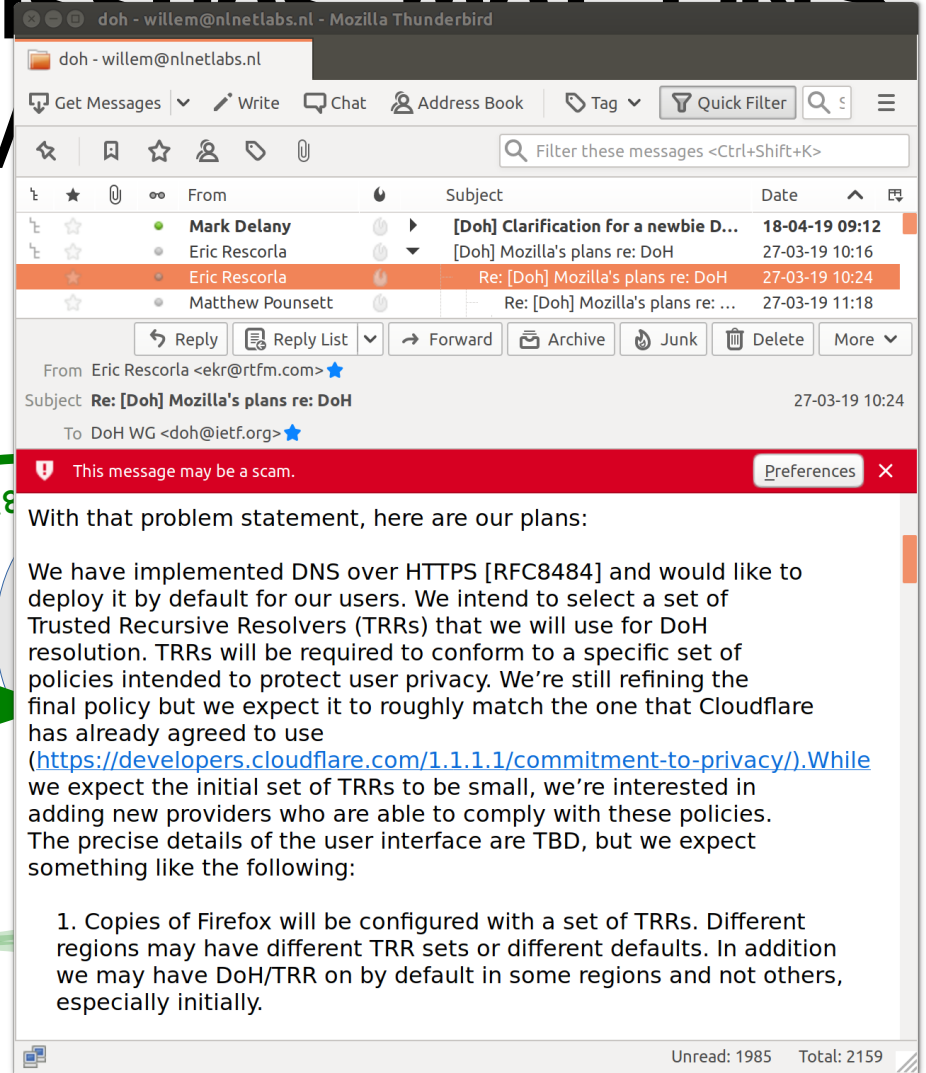
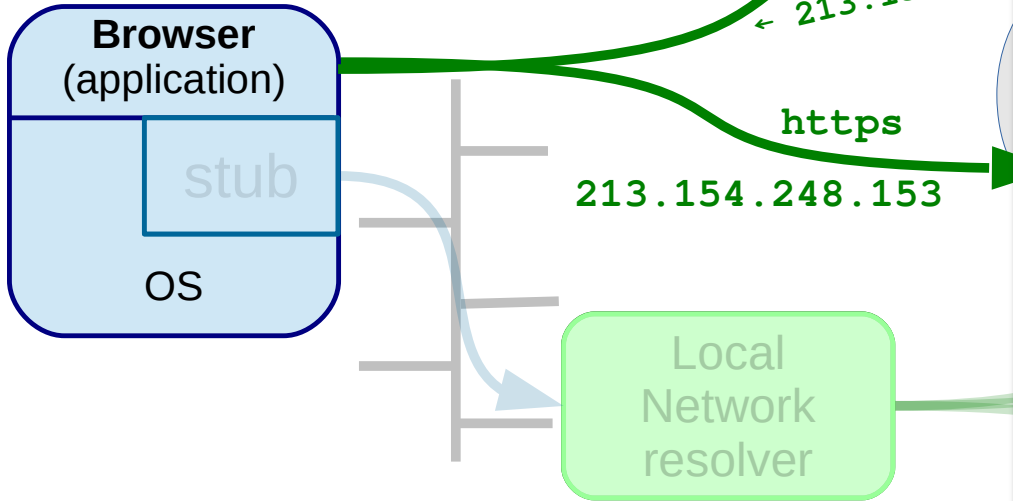


Overall Encryptie

Privacy issues met DNS DNS over

- RFC8484

- + Onmogelijk te detecteren/blokkeren

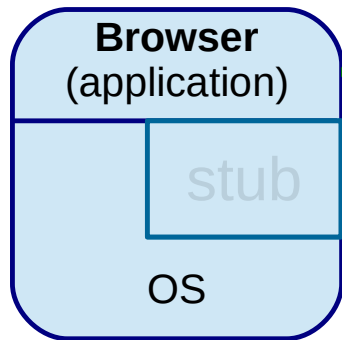


Overall
Encryptie

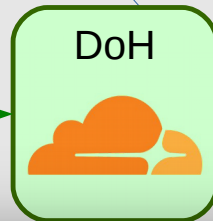
Privacy issues met DNS DNS over HTTPS (DoH)

- RFC8484

- + Onmogelijk te detecteren/blokkeren



snow.nl A →
← 213.154.248.153
https
213.154.248.153

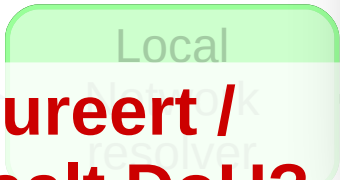


2. PRINCIPLES

Within this guiding principle, we identify two more specific principles:

- Modularize the design along tussle boundaries, so that one tussle does not spill over and distort unrelated issues.
- Design for choice, to permit the different players to express their preferences.

- **Wie stuurt / configureert / gebruikt / bepaalt DoH?**





D'OH...

The Simpsons™ © 2004 Fox



NUTS!



DONUTS

© 2004 Fox

Mmm...